



MAX PLANCK INSTITUTE  
FOR DYNAMICS OF COMPLEX  
TECHNICAL SYSTEMS  
MAGDEBURG



COMPUTATIONAL METHODS IN  
SYSTEMS AND CONTROL THEORY

# Low-rank Tensor Methods for Optimal Control of Uncertain Flow Problems

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Joint work with  
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MS PDE-constrained optimization under uncertainty  
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Supported by:



## 1. Introduction

Big data in flow control

Low-rank solvers for high-dimensional problems

PDE-constrained optimization under uncertainty

## 2. Optimal Control of Unsteady Navier-Stokes Equations under Uncertainty

## 3. Conclusions



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- Why not starting with a compressed data format from the very beginning, and computing all trajectories **all-at-once**?
- This is the idea we pursue in this talk, where we store the data in a compressed (low-rank)  $1 + 10 + 1(4)$ -way tensor! (Above example  $\rightsquigarrow$  terabyte-range.)

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- data are unpredictable, e.g., wind shear ("**epistemic uncertainty**").

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[BELLMAN '57]

Increase in matrix size of discretized differential operator for  $h \rightarrow \frac{h}{2}$  by factor  $2^d$ .

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$$(I \otimes A + A \otimes I)x =: \mathcal{A}x = b \quad \iff \quad AX + XA^T = B$$

with  $x = \text{vec}(X)$  and  $b = \text{vec}(B)$  with low-rank right hand side  $B \approx b_1 b_2^T$ .

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- Hence,  $\mathcal{A}\text{vec}(X_k) = \mathcal{A}\text{vec}(V_k W_k^T) = \text{vec} \left( [AV_k, V_k][W_k, AW_k]^T \right)$

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- Hence,  $\mathcal{A}\text{vec}(X_k) = \mathcal{A}\text{vec}(V_k W_k^T) = \text{vec}([AV_k, V_k][W_k, AW_k]^T)$
- The rank of  $[AV_k \ V_k] \in \mathbb{R}^{n, 2r}$ ,  $[W_k \ AW_k] \in \mathbb{R}^{n_t, 2r}$  increases but can be controlled using truncation.  $\rightsquigarrow$  **Low-rank Krylov subspace solvers**.

[KRESSNER/TOBLER, B/BREITEN, SAVOSTYANOV/DOLGOV, ...].

We consider the problem:

$$\min_{y \in \mathcal{Y}, u \in \mathcal{U}} \mathcal{J}(y, u) \quad \text{subject to} \quad c(y, u) = 0,$$

where we assume that

- $c(y, u) = 0$  represents a (linear or nonlinear) **PDE** (system) **with uncertain coefficient(s)**;
- the state  $y$  and control  $u$  are random fields, related by a sufficiently **smooth map**  $y = \mathcal{S}(u)$ ;
- the **cost functional**  $\mathcal{J}$  is a real-valued Fréchet-differentiable functional on  $\mathcal{Y} \times \mathcal{U}$ .

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Using low-rank tensor techniques, we need  $\approx 7 \cdot 10^7$  **bytes** = 70 **GB** to solve the KKT system in MATLAB in less than one hour!

## 1. Introduction

## 2. Optimal Control of Unsteady Navier-Stokes Equations under Uncertainty

- Model problems

- Numerical discretization techniques

- The tensor train format

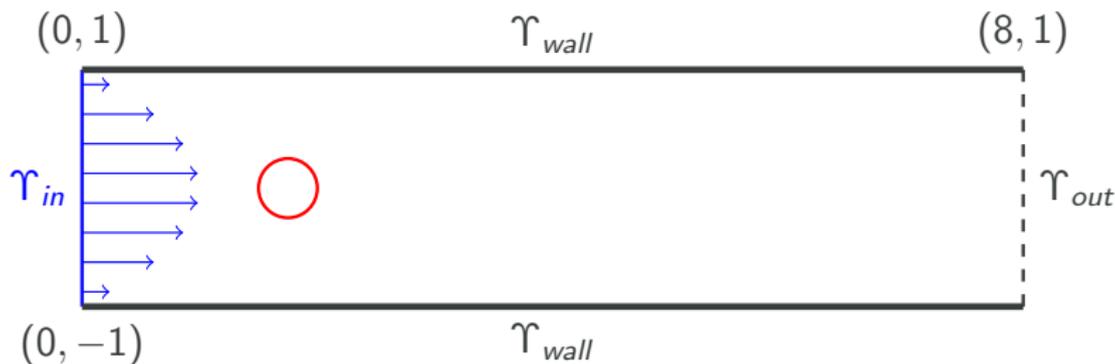
- Alternating linear solvers

- Numerical Experiments

## 3. Conclusions

## Model Problem 1: uncertain flow past a circular obstacle

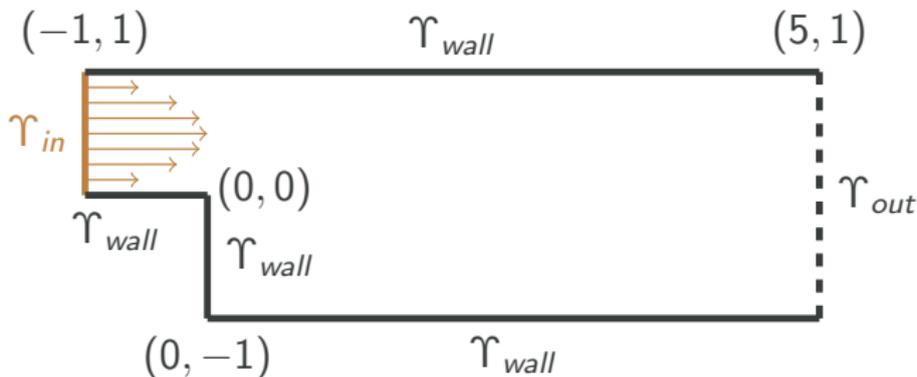
▶ to Numerical Experiments



- We model this as a **boundary control problem**.
- Our constraint  $c(y, u) = 0$  is given by the unsteady incompressible Navier-Stokes equations with **uncertain viscosity**  $\nu := \nu(\omega) /$  **inflow condition**  $\theta_2(t, x, \omega) = 0$  and

$$\theta_1(t, x, \omega) = \left( (1 + x_2)(1 - x_2) + \sum_{k=1}^m k^{-\gamma-1/2} \cdot \sin(\pi k x_2) \cdot \xi_k(\omega) \right) (1 - e^{-t}).$$

## Model Problem 2: uncertain flow in a backward facing step domain



- We model this as a **boundary control problem**.
- Our constraint  $c(y, u) = 0$  is given by the unsteady incompressible Navier-Stokes equations with **uncertain inflow condition**  $\theta_2(t, x, \omega) = 0$  and

$$\theta_1(t, x, \omega) = \left( 4x_2(1 - x_2) + \frac{1}{2} \sum_{k=1}^m k^{-\gamma-1/2} \sin(2\pi kx_2) \xi_k(\omega) \right) (1 - e^{-t}).$$

Minimize:

$$\mathcal{J}(v, u) = \frac{1}{2} \|\operatorname{curl} v\|_{L^2(0, T; \mathcal{D}) \otimes L^2(\Omega)}^2 + \frac{\beta}{2} \|u\|_{L^2(0, T; \mathcal{D}) \otimes L^2(\Omega)}^2 \quad (1)$$

subject to

$$\begin{aligned} \frac{\partial v}{\partial t} - \nu \Delta v + (v \cdot \nabla) v + \nabla p &= 0, & \text{in } \mathcal{D}, \\ -\nabla \cdot v &= 0, & \text{in } \mathcal{D}, \\ v &= \theta, & \text{on } \Upsilon_{in}, \\ v &= 0, & \text{on } \Upsilon_{wall}, \\ \frac{\partial v}{\partial n} &= u, & \text{on } \Upsilon_c, \\ \frac{\partial v}{\partial n} &= 0, & \text{on } \Upsilon_{out}, \\ v(\cdot, 0, \cdot) &= v_0, & \text{in } \mathcal{D}. \end{aligned} \quad (2)$$

We assume

- $\nu(\omega) = \nu_0 + \nu_1 \xi(\omega)$ ,  $\nu_0, \nu_1 \in \mathbb{R}^+$ ,  $\xi \sim \mathcal{U}(-1, 1)$ .
- $\mathbb{P}(\omega \in \Omega : \nu(\omega) \in [\nu_{\min}, \nu_{\max}]) = 1$ , for some  $0 < \nu_{\min} < \nu_{\max} < +\infty$ .
- $\Rightarrow$  velocity  $v$ , control  $u$  and pressure  $p$  are random fields on  $L^2(\Omega)$ .
- $L^2(\Omega) := L^2(\Omega, \mathcal{F}, \mathbb{P})$  is a complete probability space.
- $L^2(0, T; \mathcal{D}) := L^2(\mathcal{D}) \times L^2(\mathcal{T})$ .

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## Computational challenges

- Nonlinearity (due to the nonlinear convection term  $(v \cdot \nabla)v$ ).
- Uncertainty (due to random  $\nu(\omega)$  and  $\sigma(t, x, \omega)$ ).
- High dimensionality (of the resulting linear/optimality systems).

OTD Strategy and Picard (Oseen) Iteration  $\rightsquigarrow$ 

## state equation

$$\begin{aligned}
 v_t - \nu \Delta v + (\bar{v} \cdot \nabla) v + \nabla p &= 0 \\
 \nabla \cdot v &= 0 + \text{boundary conditions}
 \end{aligned}$$

## adjoint equation

$$\begin{aligned}
 -\chi_t - \Delta \chi - (\bar{v} \cdot \nabla) \chi + (\nabla \bar{v})^T \chi + \nabla \mu &= -\text{curl}^2 v \\
 \nabla \cdot \chi &= 0 \\
 \text{on } \Upsilon_{wall} \cup \Upsilon_{in} : \quad \chi &= 0 \\
 \text{on } \Upsilon_{out} \cup \Upsilon_c : \quad \frac{\partial \chi}{\partial n} &= 0 \\
 \chi(\cdot, T, \cdot) &= 0
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## gradient equation

$$\beta u + \chi|_{\Upsilon_c} = 0.$$

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- $\bar{v}$  denotes the velocity from the previous Oseen iteration.
- Having solved this system, we update  $\bar{v} = v$  until convergence.

- Velocity  $v$  and control  $u$  are of the form

$$z(t, x, \omega) = \sum_{k=0}^{P-1} \sum_{j=1}^{J_v} z_{jk}(t) \phi_j(x) \psi_k(\xi) = \sum_{k=0}^{P-1} z_k(t, x) \psi_k(\xi).$$

- Pressure  $p$  is of the form

$$p(t, x, \omega) = \sum_{k=0}^{P-1} \sum_{j=1}^{J_p} p_{jk}(t) \tilde{\phi}_j(x) \psi_k(\xi) = \sum_{k=0}^{P-1} p_k(t, x) \psi_k(\xi).$$

- Here,

- $\{\phi_j\}_{j=1}^{J_v}$  and  $\{\tilde{\phi}_j\}_{j=1}^{J_p}$  are Q2–Q1 finite elements (inf-sup stable);
- $\{\psi_k\}_{k=0}^{P-1}$  are Legendre polynomials.

- Implicit Euler/dG(0) used for temporal discretization.

Linearization and SGFEM discretization yields the following saddle point system

$$\underbrace{\begin{bmatrix} M_y & 0 & L^* \\ 0 & M_u & N^T \\ L & N & 0 \end{bmatrix}}_A \underbrace{\begin{bmatrix} y \\ u \\ \lambda \end{bmatrix}}_x = \underbrace{\begin{bmatrix} f \\ 0 \\ g \end{bmatrix}}_b.$$

Each of the block matrices in  $A$  is of the form

$$\sum_{\alpha=1}^R X_{\alpha} \otimes Y_{\alpha} \otimes Z_{\alpha},$$

corresponding to temporal, stochastic, and spatial discretizations.

Linearization and SGFEM discretization yields the following saddle point system

$$\underbrace{\begin{bmatrix} M_y & 0 & L^* \\ 0 & M_u & N^T \\ L & N & 0 \end{bmatrix}}_A \underbrace{\begin{bmatrix} y \\ u \\ \lambda \end{bmatrix}}_x = \underbrace{\begin{bmatrix} f \\ 0 \\ g \end{bmatrix}}_b.$$

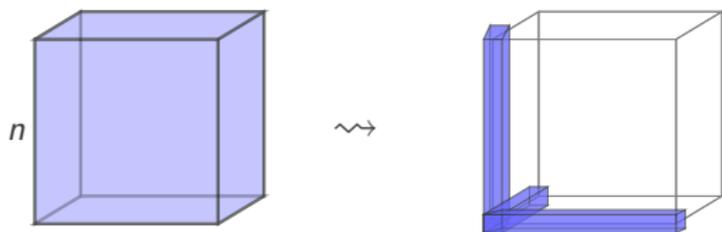
Each of the block matrices in  $A$  is of the form

$$\sum_{\alpha=1}^R X_{\alpha} \otimes Y_{\alpha} \otimes Z_{\alpha},$$

corresponding to temporal, stochastic, and spatial discretizations.

Size:  $\sim 3n_t P(J_v + J_p)$ , e.g., for  $P = 10$ ,  $n_t = 2^{10}$ ,  $J \approx 10^5 \rightsquigarrow \approx 10^9$  unknowns!

## Separation of variables and low-rank approximation



- Approximate:  $\underbrace{\mathbf{x}(i_1, \dots, i_d)}_{\text{tensor}} \approx \underbrace{\sum_{\alpha} \mathbf{x}_{\alpha}^{(1)}(i_1) \mathbf{x}_{\alpha}^{(2)}(i_2) \cdots \mathbf{x}_{\alpha}^{(d)}(i_d)}_{\text{tensor product decomposition}}.$

Goals:

- Store and manipulate  $x$
- Solve equations  $Ax = b$

$\mathcal{O}(dn)$  cost instead of  $\mathcal{O}(n^d)$ .

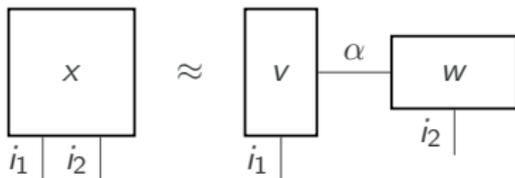
$\mathcal{O}(dn^2)$  cost instead of  $\mathcal{O}(n^{2d})$ .



- Discrete *separation of variables*:

$$\begin{bmatrix} x_{1,1} & \cdots & x_{1,n} \\ \vdots & & \vdots \\ x_{n,1} & \cdots & x_{n,n} \end{bmatrix} = \sum_{\alpha=1}^r \begin{bmatrix} v_{1,\alpha} \\ \vdots \\ v_{n,\alpha} \end{bmatrix} [w_{\alpha,1} \quad \cdots \quad w_{\alpha,n}] + \mathcal{O}(\varepsilon).$$

- Diagrams:



- Rank  $r \ll n$ .
- $\text{mem}(v) + \text{mem}(w) = 2nr \ll n^2 = \text{mem}(x)$ .
- Singular Value Decomposition (SVD)*  
 $\implies \varepsilon(r)$  optimal w.r.t. spectral/Frobenius norm.

## Tensor Trains/Matrix Product States

[WILSON '75, WHITE '93, VERSTRAETE '04, OSELEDETS '09/'11]

For indices

$$\overline{i_p \dots i_q} = (i_p - 1)n_{p+1} \dots n_q + (i_{p+1} - 1)n_{p+2} \dots n_q + \dots + (i_{q-1} - 1)n_q + i_q,$$

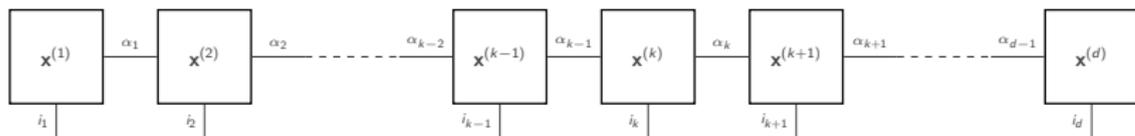
the TT format can be expressed as

$$x(\overline{i_1 \dots i_d}) = \sum_{\alpha=1}^r \mathbf{x}_{\alpha_1}^{(1)}(i_1) \cdot \mathbf{x}_{\alpha_1, \alpha_2}^{(2)}(i_2) \cdot \mathbf{x}_{\alpha_2, \alpha_3}^{(3)}(i_3) \dots \mathbf{x}_{\alpha_{d-1}, \alpha_d}^{(d)}(i_d)$$

or

$$x(\overline{i_1 \dots i_d}) = \mathbf{x}^{(1)}(i_1) \dots \mathbf{x}^{(d)}(i_d), \quad \mathbf{x}^{(k)}(i_k) \in \mathbb{R}^{r_{k-1} \times r_k} \text{ w/ } r_0, r_d = 1,$$

or



**Storage:**  $\mathcal{O}(dnr^2)$  instead of  $\mathcal{O}(n^d)$ .

Always work with *factors*  $\mathbf{x}^{(k)} \in \mathbb{R}^{r_{k-1} \times n_k \times r_k}$  instead of **full tensors**.

- Sum  $z = x + y \rightsquigarrow$  increase of tensor rank  $r_z = r_x + r_y$ .
- TT format for a high-dimensional operator

$$A(\overline{i_1 \dots i_d}, \overline{j_1 \dots j_d}) = \mathbf{A}^{(1)}(i_1, j_1) \cdots \mathbf{A}^{(d)}(i_d, j_d)$$

- *Matrix-vector* multiplication  $y = Ax$ ;  $\rightsquigarrow$  tensor rank  $r_y = r_A \cdot r_x$ .
- Additions and multiplications *increase* TT ranks.
- *Decrease* ranks quasi-optimally via QR and SVD.



# Solving KKT System using TT Format

The dimensionality of the saddle point system is vast  $\Rightarrow$  use **tensor structure** and low tensor ranks.

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Use tensor train format to approximate the solution as

$$\mathbf{y}(i_1, \dots, i_d) \approx \sum_{\alpha_1 \dots \alpha_{d-1}=1}^{r_1 \dots r_{d-1}} \mathbf{y}_{\alpha_1}^{(1)}(i_1) \mathbf{y}_{\alpha_1, \alpha_2}^{(2)}(i_2) \cdots \mathbf{y}_{\alpha_{d-2}, \alpha_{d-1}}^{(d-1)}(i_{d-1}) \mathbf{y}_{\alpha_{d-1}}^{(d)}(i_d),$$

and represent the coefficient matrix as

$$\mathcal{A}(i_1 \cdots i_d, j_1 \cdots j_d) \approx \sum_{\beta_1 \dots \beta_{d-1}=1}^{R_1 \dots R_{d-1}} \mathbf{A}_{\beta_1}^{(1)}(i_1, j_1) \mathbf{A}_{\beta_1, \beta_2}^{(2)}(i_2, j_2) \cdots \mathbf{A}_{\beta_{d-1}}^{(d)}(i_d, j_d),$$

where the multi-index  $\mathbf{i} = (i_1, \dots, i_d)$  is implied by the parametrization of the approximate solutions of the form

$$\mathbf{z}(t, \xi_1, \dots, \xi_N, \mathbf{x}), \quad \mathbf{z} = \mathbf{y}, \mathbf{u}, \mathbf{p},$$

i.e., **solution vectors are represented by  $d$ -way tensor with  $d = N + 2$ .**



## Central Question

How to solve  $Ax = b$  if Krylov solvers become too expensive?



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Data are given in TT format:

- $A(i, j) = \mathbf{A}^{(1)}(i_1, j_1) \cdots \mathbf{A}^{(d)}(i_d, j_d)$ .
- $b(i) = \mathbf{b}^{(1)}(i_1) \cdots \mathbf{b}^{(d)}(i_d)$ .

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Use a new block-variant of *Alternating Least Squares* in a new block TT format to overcome difficulties with indefiniteness of KKT system matrix.

- If  $A = A^\top > 0$ : minimize  $J(x) = x^\top Ax - 2x^\top b$ .

### *Alternating Least Squares (ALS):*

- replace  $\min_x J(x)$  by iteration
- for  $k = 1, \dots, d$ ,  
 solve  $\min_{\mathbf{x}^{(k)}} J(\mathbf{x}^{(1)}(i_1) \cdots \mathbf{x}^{(k)}(i_k) \cdots \mathbf{x}^{(d)}(i_d))$ .  
 (all other blocks are fixed)

size  $n^d$

size  $r^2 n$



1.  $\hat{\mathbf{x}}^{(1)} = \arg \min_{\mathbf{x}^{(1)}} J(\mathbf{x}^{(1)}(i_1)\mathbf{x}^{(2)}(i_2)\mathbf{x}^{(3)}(i_3))$



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5. repeat 1.–4. until convergence



If we differentiate  $J$  w.r.t. TT blocks, we see that...

- ... each step means solving a *Galerkin linear system*

$$\left( X_{\neq k}^\top A X_{\neq k} \right) \hat{\mathbf{x}}^{(k)} = \left( X_{\neq k}^\top \mathbf{b} \right) \in \mathbb{R}^{nr^2}.$$

- $$X_{\neq k} = \underbrace{\text{TT} \left( \hat{\mathbf{x}}^{(1)} \dots \hat{\mathbf{x}}^{(k-1)} \right)}_{n^{k-1} \times r_{k-1}} \otimes \underbrace{I}_{n \times n} \otimes \underbrace{\text{TT} \left( \mathbf{x}^{(k+1)} \dots \mathbf{x}^{(d)} \right)}_{n^{d-k} \times r_k}.$$

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## Properties of ALS include:

- + Effectively 1D complexity in a prescribed format.
- Tensor format (ranks) is fixed and cannot be adapted.
- Convergence may be very slow, stagnation is likely.



- Density Matrix Renormalization Group (DMRG) [WHITE '92]  
– updates *two* blocks  $\mathbf{x}^{(k)}\mathbf{x}^{(k+1)}$  *simultaneously*.
- Alternating Minimal Energy (AMEn) [DOLGOV/SAVOSTYANOV '13]  
– *augments*  $X_{\neq k}$  by a TT block of the *residual*  $\mathbf{z}^{(k)}$ .

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But... , what about saddle point systems  $A$ ?

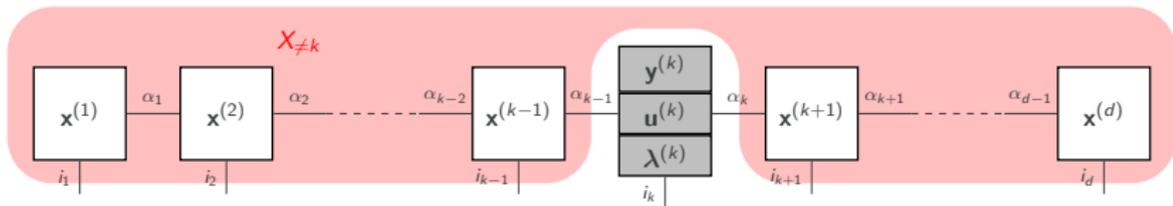
- Recall our KKT system:

$$\underbrace{\begin{bmatrix} M_y & 0 & L^* \\ 0 & M_u & N^T \\ L & N & 0 \end{bmatrix}}_A \begin{bmatrix} y \\ u \\ \lambda \end{bmatrix} = \begin{bmatrix} f \\ 0 \\ g \end{bmatrix}.$$

- The whole matrix is **indefinite**  $\Rightarrow X_{\neq k}^T A X_{\neq k}$  can be degenerate.

- Work-around: Block TT representation

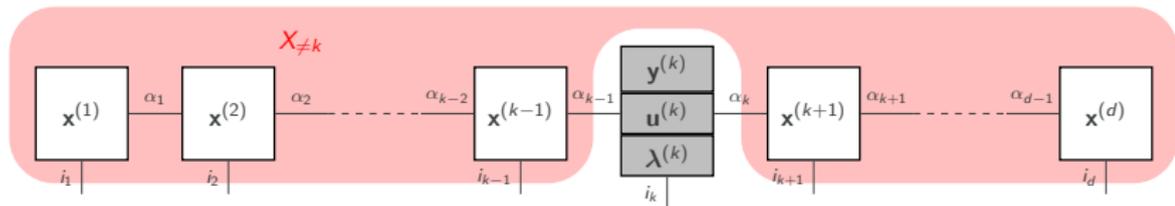
$$\begin{bmatrix} y \\ u \\ \lambda \end{bmatrix} = \mathbf{x}_{\alpha_1}^{(1)} \otimes \dots \otimes \begin{bmatrix} \mathbf{y}_{\alpha_{k-1}, \alpha_k}^{(k)} \\ \mathbf{u}_{\alpha_{k-1}, \alpha_k}^{(k)} \\ \boldsymbol{\lambda}_{\alpha_{k-1}, \alpha_k}^{(k)} \end{bmatrix} \otimes \dots \otimes \mathbf{x}_{\alpha_{d-1}}^{(d)}.$$



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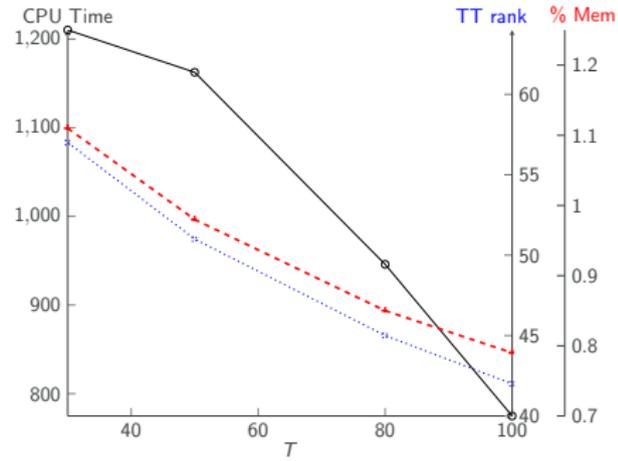
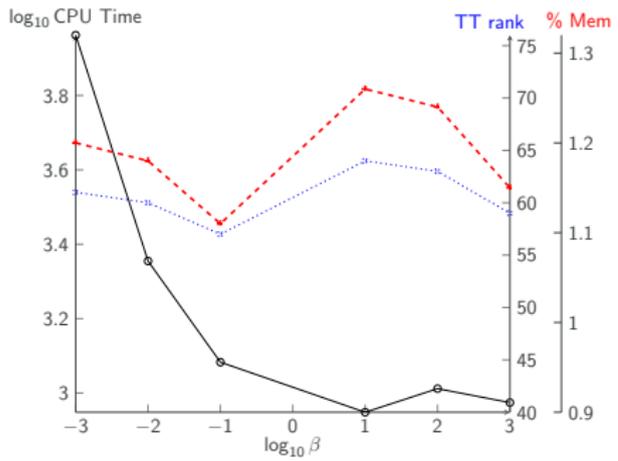
- $X_{\neq k}$  is the *same* for  $y, u, \lambda$ .
- Project each *submatrix*:

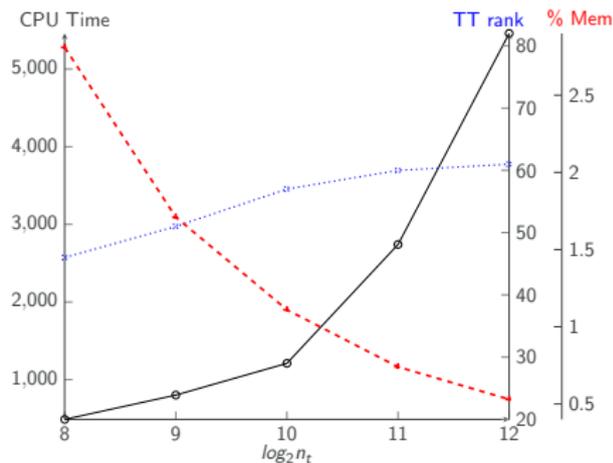
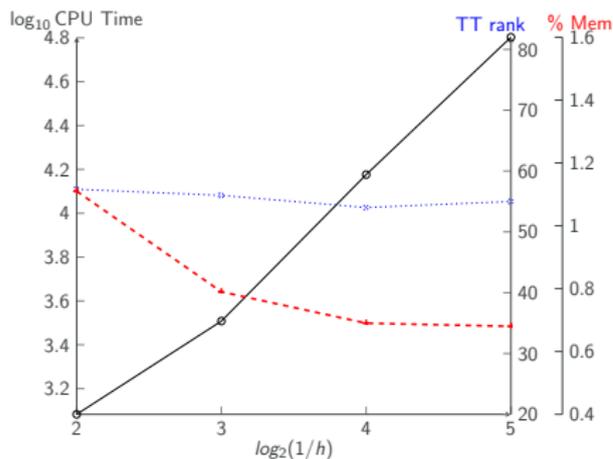
$$\begin{bmatrix} \hat{M}_y & 0 & \hat{L}^* \\ 0 & \hat{M}_u & \hat{N}^\top \\ \hat{L} & \hat{N} & 0 \end{bmatrix} \begin{bmatrix} y^{(k)} \\ u^{(k)} \\ \lambda^{(k)} \end{bmatrix} = \begin{bmatrix} \hat{f} \\ 0 \\ \hat{g} \end{bmatrix}, \quad \widehat{(\cdot)} = X_{\neq k}^\top (\cdot) X_{\neq k}.$$

## Vary one of the default parameters:

- TT truncation tolerance  $\varepsilon = 10^{-4}$ ,
- mean viscosity  $\nu_0 = 1/20$ ,
- uncertainty  $\nu_1 = 1/80$ ,
- regularization/penalty parameter  $\beta = 10^{-1}$ ,
- number of time steps:  $n_t = 2^{10}$ ,
- time horizon  $T = 30$ ,
- spatial grid size  $h = 1/4 \rightsquigarrow J = 2488$ ,
- max. degree of Legendre polynomials:  $P = 8$ .

Solve projected linear systems using block-preconditioned GMRES using efficient approximation of Schur complement [B/ONWUNTA/STOLL 2016].





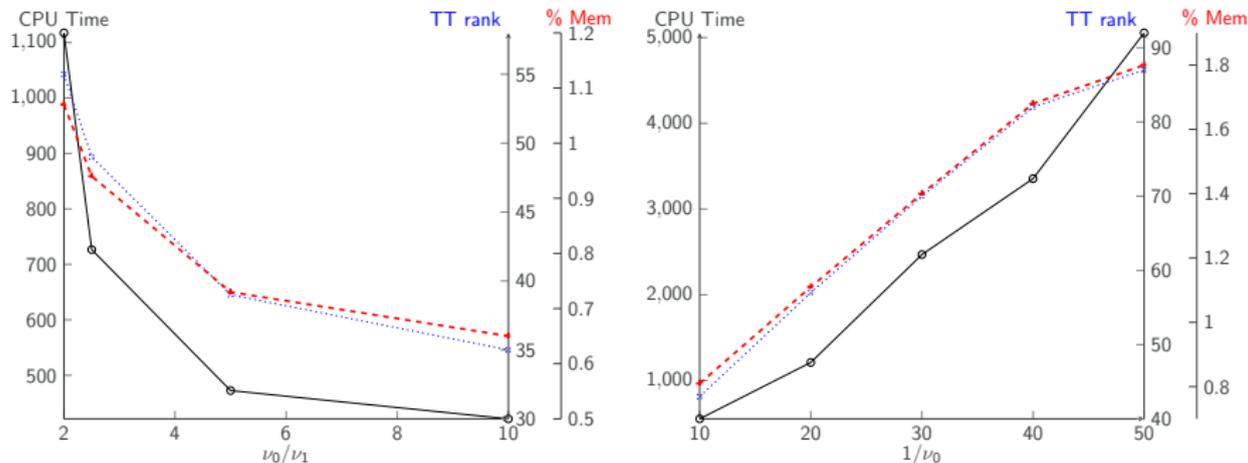
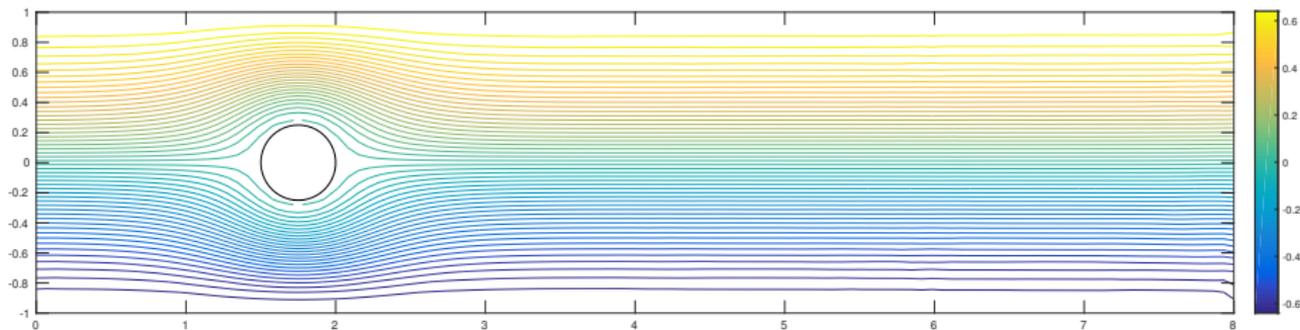
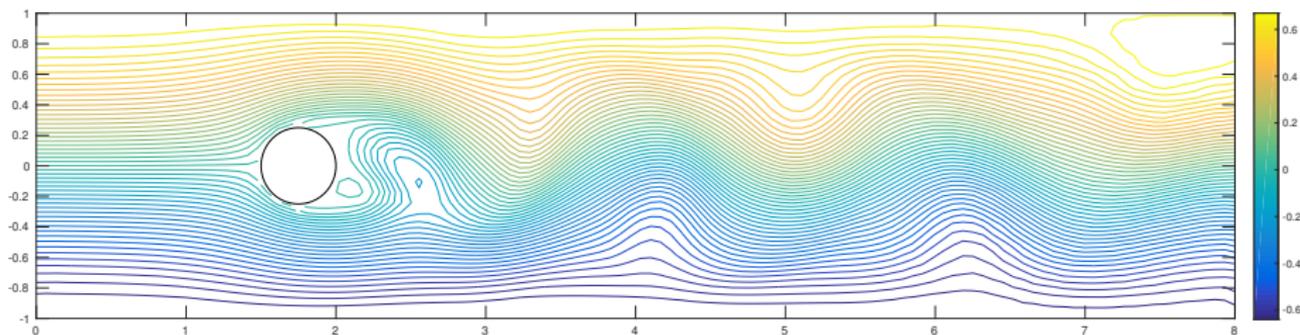
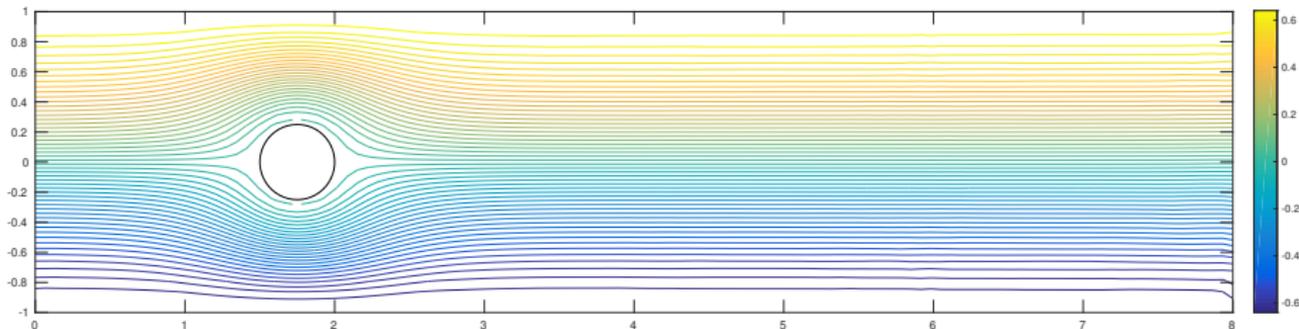
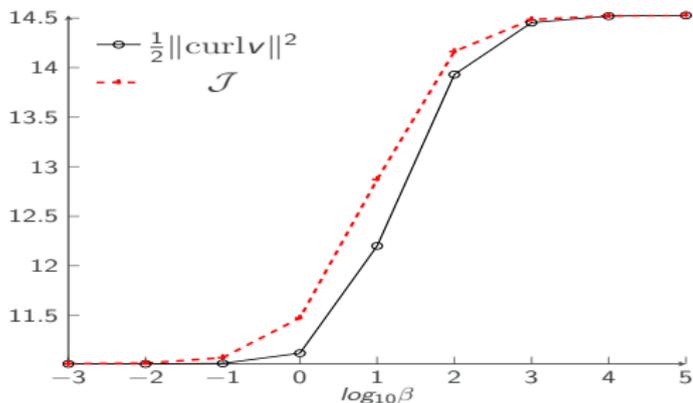


Figure: Left:  $\nu_0 = 1/10$ ,  $\nu_1$  is varied. Right:  $\nu_1$  and  $\nu_0$  are varied together as  $\nu_1 = 0.25\nu_0$

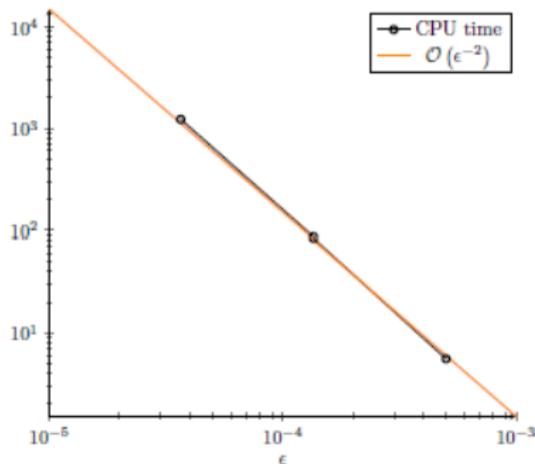




Balance all errors w.r.t.

- time, spatial, and stochastic discretization;
- TT truncation.

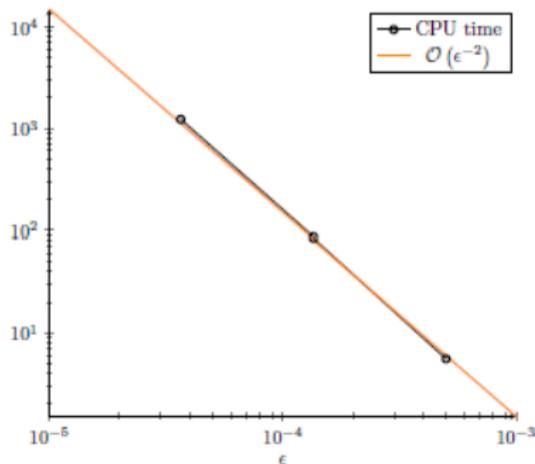
Call the **balanced error**  $\varepsilon$ .



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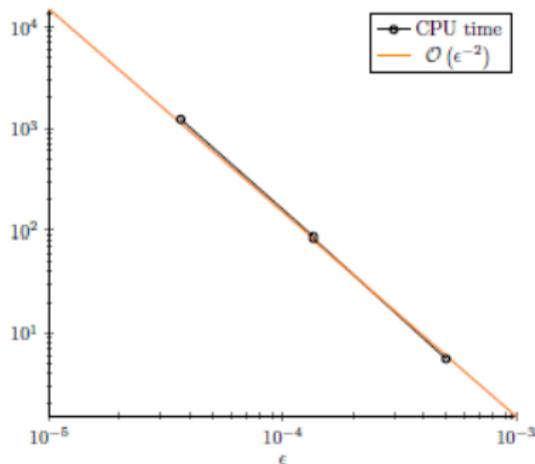


This indicates **asymptotic complexity**  $\varepsilon^{-2}$ , asymptotically equal complexity as for deterministic problem.

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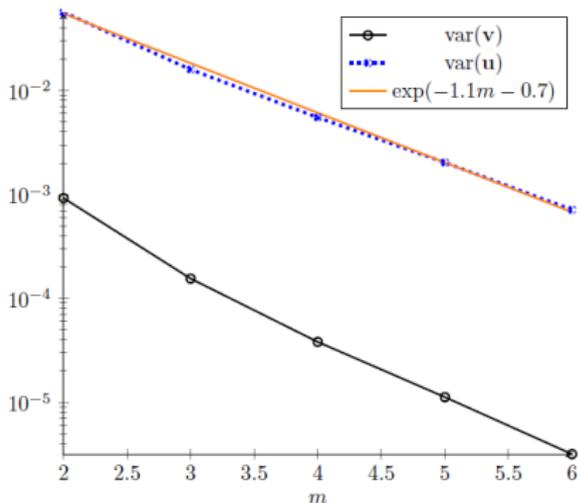
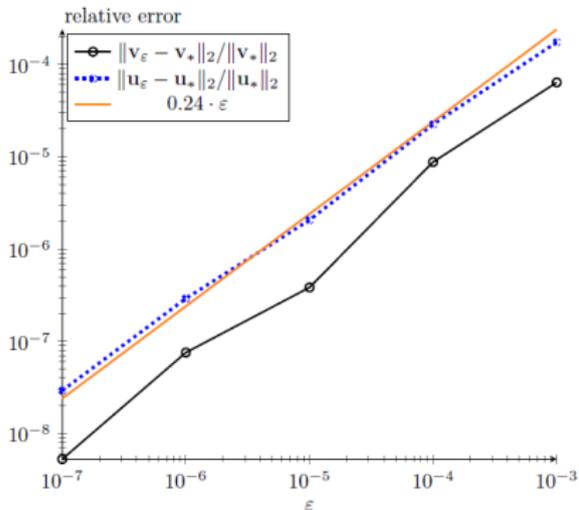
This compares favorably to

- Monte-Carlo  $\mathcal{O}(\varepsilon^{-4})$ ,
- quasi Monte-Carlo / stochastic collocation  $\mathcal{O}(\varepsilon^{-3})$ .

Relative errors for different TT approximation thresholds (left) and varying  $m$  (right)...

## Parameters:

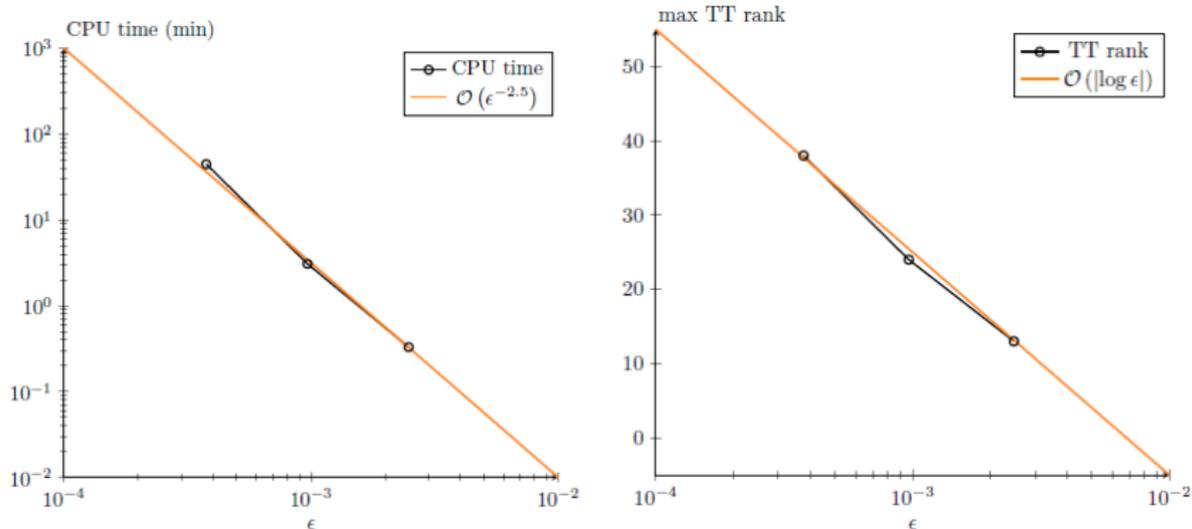
viscosity  $\nu \equiv \nu_0 = 10^{-3}$ , final time  $T = 20$ , regularization parameter  $\beta = 10^{-2}$ , KL decay rate  $\gamma = 2.5$  for inflow condition.



Total CPU time (left) and TT rank (right) versus total approximation error.

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1. Introduction
2. Optimal Control of Unsteady Navier-Stokes Equations under Uncertainty
3. Conclusions

- Low-rank tensor solver for unsteady heat and Navier-Stokes equations with uncertain viscosity.
- Similar techniques already used for Stokes(-Brinkman) optimal control problems.
- Adapted AMEn (TT) solver to saddle point systems.
- With 1 stochastic parameter, the scheme reduces complexity by up to 2–3 orders of magnitude.
- To consider next:

- Low-rank tensor solver for unsteady heat and Navier-Stokes equations with uncertain viscosity.
- Similar techniques already used for Stokes(-Brinkman) optimal control problems.
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- To consider next:
  - many parameters coming from uncertain geometry or Karhunen-Loève expansion of random fields;  
Basic observation: the more parameters, the more significant is the complexity reduction w.r.t. memory — up to a factor of  $10^9$  for the control problem for a backward facing step set-up.
  - HPC implementation of AMEn-like solver to deal with even larger problems.



P. Benner, S. Dolgov, A. Onwunta, and M. Stoll.

Low-rank solvers for unsteady Stokes-Brinkman optimal control problem with random data.

COMPUTER METHODS IN APPLIED MECHANICS AND ENGINEERING, 304:26–54, 2016.



P. Benner, A. Onwunta, and M. Stoll.

Low rank solution of unsteady diffusion equations with stochastic coefficients.

SIAM/ASA JOURNAL ON UNCERTAINTY QUANTIFICATION, 3(1):622–649, 2015.



P. Benner, A. Onwunta, and M. Stoll.

Block-diagonal preconditioning for optimal control problems constrained by PDEs with uncertain inputs.

SIAM JOURNAL ON MATRIX ANALYSIS AND APPLICATIONS, 37(2):491–518, 2016.



P. Benner, S. Dolgov, A. Onwunta, and M. Stoll.

Solving optimal control problems governed by random Navier-Stokes equations using low-rank methods. [arXiv:1703.06097](https://arxiv.org/abs/1703.06097), March 2017.