

A survey about moment matching model order reduction in computational electromagnetism

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- Four industrial partners. (CST, X-FAB, ...)



Maxwell's equations

$$\begin{pmatrix} M_{\epsilon} & 0\\ 0 & M_{\mu} \end{pmatrix} \dot{x}(t) = \begin{pmatrix} -M_{\sigma} & C_{H}\\ -C_{E} & 0 \end{pmatrix} x(t) + \mathcal{B}u(t), \ y(t) = \mathcal{C}x(t)$$

- Block structure with respect to electric and magnetic field strength.
- Discretization leads to high-dimensional model problem.
- For model order reduction.
 - Structure-preserving reduced order model.
 - · Efficient solution technique for shifted linear systems.



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Outline

- Moment matching methods in model order reduction
- Efficient offline-stage of moment matching methods
 - Modified adaptive-order rational Arnoldi method
 - Recycling Krylov subspace methods
- Numerical experiments
- Conclusion



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Moment matching in a nutshell

Taylor expansion Let $\mathcal{H}(s) = \sum_{j=0}^{\infty} C X^{(j)}(s_i)(s-s_i)^j$ denote the transfer function at $s_i \in \mathbb{C}$, where $X^{(j)}(s_i) = \left[-(s_i \mathcal{E} - \mathcal{A})^{-1} \mathcal{E}\right]^{j-1} (s_i \mathcal{E} - \mathcal{A})^{-1} \mathcal{B}.$



Moment matching in a nutshell





Moment matching in a nutshell



Padé approximation

Galerkin projection leads to $\mathcal{H}^{(j)}(\mathbf{s}_i) = \tilde{\mathcal{H}}^{(j)}(\mathbf{s}_i)$ for all $j = 0, \dots, j_i - 1$.



Let $\mathbf{Y}^{(j)}(\mathbf{s}_i) \equiv \mathbb{C}\mathbf{X}^{(j)}(\mathbf{s}_i)$ denote the output moments.



• Note that $r_j(s) = -(s\mathcal{E} - \mathcal{A})^{-1}\mathcal{E}r_{j-1}(s)$ for all $s \in S_j$.



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• Note that $r_j(s) = -(s\mathcal{E} - \mathcal{A})^{-1}\mathcal{E}r_{j-1}(s)$ for all $s \in S_j$.



Lemma ([B., B., 2013]) Let $\mathcal{P}(s) = s\mathcal{E} - \mathcal{A}$ and $\tilde{\mathcal{P}}(s) = s\mathcal{\tilde{E}} - \mathcal{\tilde{A}}$. $|\mathbf{Y}^{(0)}(s) - \mathbf{\tilde{Y}}^{(0)}(s)| \leq |\mathcal{C}\mathcal{P}(s)^{-1}| \cdot |\mathbf{h}_n(s)| > 1000$



Lemma ([B., B., 2013])
Let
$$\mathcal{P}(s) = s\mathcal{E} - \mathcal{A}$$
 and $\tilde{\mathcal{P}}(s) = s\mathcal{\tilde{E}} - \mathcal{\tilde{A}}$.
 $|Y^{(0)}(s) - \tilde{Y}^{(0)}(s)| \leq |\mathcal{CP}(s)^{-1}| \cdot |h_n(s)| \approx 10$

$$h_n(\boldsymbol{s}_{m+1}) = \mathcal{B} - \mathcal{P}(\boldsymbol{s}_{m+1}) V \tilde{\mathcal{P}}(\boldsymbol{s}_{m+1})^{-1} \tilde{\mathcal{B}}$$



Lemma ([B., B., 2013])
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 $|Y^{(0)}(s) - \tilde{Y}^{(0)}(s)| \leq |\mathbb{CP}(s)^{-1}| \cdot |h_n(s)| \simeq 1$

$$h_n(s_{m+1})\mathbf{u} = \mathcal{B}\mathbf{u} - \mathcal{P}(s_{m+1})V\tilde{\mathcal{P}}(s_{m+1})^{-1}\tilde{\mathcal{B}}\mathbf{u}$$



Lemma ([B., B., 2013])
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$$h_n(\boldsymbol{s}_{m+1})\mathbf{u} = \mathcal{B}\mathbf{u} - \mathcal{P}(\boldsymbol{s}_{m+1})V\tilde{\mathcal{P}}(\boldsymbol{s}_{m+1})^{-1}\tilde{\mathcal{B}}\mathbf{u}$$

$$\mathcal{P}(\boldsymbol{s}_{m+1}) \boldsymbol{V} \tilde{\mathcal{P}}(\boldsymbol{s}_{m+1})^{-1} \tilde{\mathcal{B}} \mathbf{u} \approx \boldsymbol{\mathcal{B}} \mathbf{u}$$



Lemma ([B., B., 2013]) Let $\mathcal{P}(s) = s\mathcal{E} - \mathcal{A}$ and $\tilde{\mathcal{P}}(s) = s\mathcal{\tilde{E}} - \mathcal{\tilde{A}}$. $|Y^{(0)}(s) - \tilde{Y}^{(0)}(s)| \leq |\mathcal{CP}(s)^{-1}| \cdot |h_n(s)| \approx |\mathcal{\tilde{CP}}(s)^{-1}| \cdot |h_n(s)|,$

$$h_n(s_{m+1})\mathbf{u} = \mathcal{B}\mathbf{u} - \mathcal{P}(s_{m+1})V\tilde{\mathcal{P}}(s_{m+1})^{-1}\tilde{\mathcal{B}}\mathbf{u}$$

$$\mathcal{P}(\boldsymbol{s}_{m+1}) \boldsymbol{V} \tilde{\mathcal{P}}(\boldsymbol{s}_{m+1})^{-1} \tilde{\mathcal{B}} \mathbf{u} ~\approx~ \mathcal{B} \mathbf{u}$$

AORA-RK method

Determine $s_{i+1} \in \mathbb{C}$ such that

$$s_{i+1} = \arg\max_{s\in\mathbb{S}} |\tilde{\mathbb{C}}\tilde{\mathbb{P}}(s)^{-1}| \cdot |h_n(s)| \text{ with } \mathbb{S} \subset \iota[f_{\min}, f_{\max}].$$



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Simplified example for two expansion points $s_1 \in \mathbb{C}$ and $s_2 \in \mathbb{C}$.



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Simplified example for two expansion points $s_1 \in \mathbb{C}$ and $s_2 \in \mathbb{C}$.

 $S_1 = \{\boldsymbol{s}_1\}$



Shifted linear system:
$$(s_1 \mathcal{E} - \mathcal{A})r_j(s_1) = r_{j-1}(s_1)$$
.



Simplified example for two expansion points $s_1 \in \mathbb{C}$ and $s_2 \in \mathbb{C}$.





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Simplified example for two expansion points $s_1 \in \mathbb{C}$ and $s_2 \in \mathbb{C}$.

ç	$s_1 = \{s_1\}$	
	$v_1(s_1)$	
	$v_2(s_1)$	
	÷	
	$v_{n_r}(s_1)$	

h



$$(\boldsymbol{s}_1 \mathcal{E} - \mathcal{A})\boldsymbol{r}_j(\boldsymbol{s}_1) = \boldsymbol{r}_{j-1}(\boldsymbol{s}_1)$$

$$(\boldsymbol{s}_2 \mathcal{E} - \mathcal{A})\boldsymbol{r}_j(\boldsymbol{s}_2) = \boldsymbol{r}_{j-1}(\boldsymbol{s}_2)$$

Expensive approach due to repeated solution of linear systems.



Simplified example for two expansion points $s_1 \in \mathbb{C}$ and $s_2 \in \mathbb{C}$.

 $v_{n_r}(s_2)$

ĉ	$B_1 = \{s_1\}$	8 ₂	$s_2 = \{s_1, s_2\}$	2
	$v_1(s_1)$		$v_1(s_1)$	
	<i>v</i> ₂ (<i>s</i> ₁)		v ₂ (s ₂)	
	÷		<i>v</i> ₃ (<i>s</i> ₂)	
	$v_{n_r}(s_1)$		<i>v</i> ₄ (<i>s</i> ₁)	

$$(\mathbf{s}_2 \mathcal{E} - \mathcal{A})\mathbf{r}_j(\mathbf{s}_2) = \mathbf{r}_{j-1}(\mathbf{s}_2)$$



Simplified example for two expansion points $s_1 \in \mathbb{C}$ and $s_2 \in \mathbb{C}$.

$S_1 = \{\boldsymbol{s}_1\}$	Sz	$s_2 = \{s_1, s_2\}$	<u>}</u> }
<i>v</i> ₁ (<i>s</i> ₁)		<i>v</i> ₁ (<i>s</i> ₁)	
<i>v</i> ₂ (<i>s</i> ₁)		v ₂ (s ₂)	
÷		<i>v</i> ₃ (<i>s</i> ₂)	
$v_{n_r}(s_1)$		$v_4(s_1)$	
		:	
		$v_{n_r}(s_2)$	

$$(\mathbf{s}_2 \mathcal{E} - \mathcal{A})\mathbf{r}_j(\mathbf{s}_2) = \mathbf{r}_{j-1}(\mathbf{s}_2)$$



Simplified example for two expansion points $s_1 \in \mathbb{C}$ and $s_2 \in \mathbb{C}$.




Modified generic rational Arnoldi method

Simplified example for two expansion points $s_1 \in \mathbb{C}$ and $s_2 \in \mathbb{C}$.





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Modified generic rational Arnoldi method

Simplified example for two expansion points $s_1 \in \mathbb{C}$ and $s_2 \in \mathbb{C}$.



$$(\boldsymbol{s}_2 \mathcal{E} - \mathcal{A}) \boldsymbol{r}_j(\boldsymbol{s}_2) = \boldsymbol{r}_{j-1}(\boldsymbol{s}_2)$$



Modified generic rational Arnoldi method

Simplified example for two expansion points $s_1 \in \mathbb{C}$ and $s_2 \in \mathbb{C}$.

$S_1 = \{S_1\}$	Sz	$s_2 = \{s_1, s_2\}$	}
$v_1(s_1)$]	$v_1(s_1)$	
$v_2(s_1)$		v ₂ (s ₂)	
:		v ₃ (s ₂)	
$v_{n_r}(s_1)$		$v_4(s_1)$	
		÷	
		$v_{n_r}(s_2)$	

$$(\boldsymbol{s}_2 \mathcal{E} - \mathcal{A}) \boldsymbol{r}_j(\boldsymbol{s}_2) = \boldsymbol{r}_{j-1}(\boldsymbol{s}_2)$$

- Significantly smaller number of solutions to linear systems.
- Consideration of the order of previous orthonormal vector sequence.
- Extension to multiple expansion points straightforward.



















Compute sequence of reduced order models with $S_{i+1} = S_i \cup \{s_{i+1}\}, s_{i+1} \in \mathbb{C}$.







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Compute sequence of reduced order models with $S_{i+1} = S_i \cup \{s_{i+1}\}, s_{i+1} \in \mathbb{C}$.





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Compute sequence of reduced order models with $S_{i+1} = S_i \cup \{s_{i+1}\}, s_{i+1} \in \mathbb{C}$. Si S_{i+1} $r_i = -(s_{i+1}\mathcal{E} - \mathcal{A})^{-1}\mathcal{E}r_{i-1}$ $v_1(s_1)$ $V_1(S_1)$ $r_1(s_{i+1})$ $V_2(S_2)$ $v_2(s_{i+1})$ $r_2(s_{i+1})$ $V_3(S_3)$ $V_3(s_{i+1})$ $r_3(s_{i+1})$ $\operatorname{argmax}_{s}|Y^{(j)}(s) - \tilde{Y}^{(j)}(s)|$



Compute sequence of reduced order models with $S_{i+1} = S_i \cup \{s_{i+1}\}, s_{i+1} \in \mathbb{C}$.



• Additionally consider expansion points $s_l \in S_i$ with $j_{l,i+1} > j_{l,i}$.



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Schur complement with J-symmetry:

 $S_i = (s_i M_{\epsilon} + M_{\sigma}) + C_E (s_i M_{\mu})^{-1} C_H$ with $S_i^{\mathrm{T}} J = J S_i$ and J = I.



Schur complement with J-symmetry:

$$S_i = (s_i M_{\epsilon} + M_{\sigma}) + C_E (s_i M_{\mu})^{-1} C_H$$
 with $S_i^{\mathrm{T}} J = J S_i$ and $J = I$.

Krylov subspace methods

$$x_k = x_0 + \mathcal{K}_k(S_i, r_0)$$
 such that $r_k = f - S_i x_k \perp \mathcal{L}_k$

Either employ $\mathcal{L}_k \equiv \mathcal{K}_k(S_i, r_0)$ or $\mathcal{L}_k \equiv \mathcal{K}_k(S_i^{\mathrm{T}}, \tilde{r}_0)$.

SQMR method

Determine $x_k = x_0 + V_k y_k$ with minimization

$$\|\rho_0 \boldsymbol{e}_1 - \underline{T}_k \boldsymbol{y}_k\| = \min_{\boldsymbol{v}} \|\rho_0 \boldsymbol{e}_1 - \underline{T}_k \boldsymbol{y}\|.$$

Unsym. Lanczos method.

$$S_i V_k = V_{k+1} \underline{I}_k$$
$$S_i^{\mathrm{T}} W_k = W_{k+1} \underline{\tilde{I}}_k$$



Schur complement with J-symmetry:

$$S_i = (s_i M_{\epsilon} + M_{\sigma}) + C_E (s_i M_{\mu})^{-1} C_H$$
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SQMR method

Determine $x_k = x_0 + V_k y_k$ with minimization

$$\|\rho_0 \boldsymbol{e}_1 - \underline{T}_k \boldsymbol{y}_k\| = \min_{\boldsymbol{u}} \|\rho_0 \boldsymbol{e}_1 - \underline{T}_k \boldsymbol{y}\|.$$

Unsym. Lanczos method.

$$S_{i}V_{k} = V_{k+1}\underline{T}_{k}$$
$$S_{i}^{\mathrm{T}}W_{k} = W_{k+1}\underline{\widetilde{T}}_{k}$$



Schur complement with J-symmetry:

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Krylov subspace methods

$$x_k = x_0 + \mathcal{K}_k(S_i, r_0)$$
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Either employ
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Determine
$$x_k = x_0 + V_k y_k$$
 with minimization

$$\|\rho_0 \boldsymbol{e}_1 - \underline{T}_k \boldsymbol{y}_k\| = \min_{\boldsymbol{v}} \|\rho_0 \boldsymbol{e}_1 - \underline{T}_k \boldsymbol{y}\|.$$

Unsym. Lanczos method.

$$S_i V_k = V_{k+1} \underline{T}_k$$
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Schur complement with J-symmetry:

$$S_i = (s_i M_{\epsilon} + M_{\sigma}) + C_E (s_i M_{\mu})^{-1} C_H$$
 with $S_i^{\mathrm{T}} J = J S_i$ and $J = I$.

Krylov subspace methods

$$x_k = x_0 + \mathcal{K}_k(S_i, r_0)$$
 such that $r_k = f - S_i x_k \perp \mathcal{L}_k$

Either employ $\mathcal{L}_k \equiv \mathcal{K}_k(S_i, r_0)$ or $\mathcal{L}_k \equiv \mathcal{K}_k(S_i^{\mathrm{T}}, \tilde{r}_0)$.

SQMR method

Determine $x_k = x_0 + V_k y_k$ with minimization

$$\|\rho_0 \boldsymbol{e}_1 - \underline{T}_k \boldsymbol{y}_k\| = \min_{\boldsymbol{y}} \|\rho_0 \boldsymbol{e}_1 - \underline{T}_k \boldsymbol{y}\|.$$

J-symmetry:
$$S_i^{T}J = JS_i$$
, i.e. $W_k = JV_k$.

Unsym. Lanczos method.

$$S_i V_k = V_{k+1} \underline{T}_k$$
$$S_i^{\mathrm{T}} W_k = W_{k+1} \underline{\widetilde{T}}_k$$

Schur complement with J-symmetry:

$$S_i = (s_i M_{\epsilon} + M_{\sigma}) + C_E (s_i M_{\mu})^{-1} C_H$$
 with $S_i^{\mathrm{T}} J = J S_i$ and $J = I$.

Krylov subspace methods

 $x_k = x_0 + \mathcal{K}_k(S_i, r_0)$ such that $r_k = f - S_i x_k \perp \mathcal{L}_k$

Either employ $\mathcal{L}_k \equiv \mathcal{K}_k(S_i, r_0)$ or $\mathcal{L}_k \equiv \mathcal{K}_k(S_i^{\mathrm{T}}, \tilde{r}_0)$.

SQMR method Determine $x_k = x_0 + V_k y_k$ with minimization $\|\rho_0 e_1 - \underline{T}_k y_k\| = \min_y \|\rho_0 e_1 - \underline{T}_k y\|$. J-symmetry: $S_i^T J = JS_i$, i.e. $W_k = JV_k$. Inefficient for multiple solution to different shifted linear systems and multiple right hand sides.

$$S_i V_k = V_{k+1} \underline{T}_k$$

$$S_i^{\mathrm{T}} W_k = W_{k+1} \underline{\widetilde{T}}_k$$

Multiple solution to linear system $S_j x_j = f_j$ with $f_j \in \mathbb{C}^{n \times p}$ and j = 1, ..., I.



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Multiple solution to linear system $S_j x_j = f_j$ with $f_j \in \mathbb{C}^{n \times p}$ and j = 1, ..., I.

Solution update:

$$x_k = x_0 + V_k y_k + U_j z$$

- Arnoldi-/Lanczos-type method: $(I - C_j (\tilde{C}_j^{\mathrm{T}} C_j)^{-1} \tilde{C}_j^{\mathrm{T}}) S_j V_k = V_{k+1} \underline{I}_k$
- Biorthogonality condition:
 - $\left[\tilde{C}_{j}, W_{k} \right] \perp_{b} \left[C_{j}, V_{k} \right]$

Recycling subspace $C_j = S_j U_j$ and $\tilde{C}_j = S_j^{\mathrm{T}} \tilde{U}_j$



Multiple solution to linear system $S_j x_j = f_j$ with $f_j \in \mathbb{C}^{n \times p}$ and j = 1, ..., I.

Recycling Krylov subspace methods

Solution update:

$$x_k = x_0 + V_k y_k + U_j z$$

- Arnoldi-/Lanczos-type method: $(I - C_j (\tilde{C}_j^{\mathrm{T}} C_j)^{-1} \tilde{C}_j^{\mathrm{T}}) S_j V_k = V_{k+1} \underline{I}_k$
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ight]$

Recycling subspace $C_j = S_j U_j$ and $\tilde{C}_j = S_j^{\mathrm{T}} \tilde{U}_j$

Employing (Harmonic) Ritz values, it follows that $U_i = [U_{i-1}, \bar{V}_k].$



Multiple solution to linear system $S_j x_j = f_j$ with $f_j \in \mathbb{C}^{n \times p}$ and j = 1, ..., I.

Solution update:

$$x_k = x_0 + V_k y_k + U_j z$$

- Arnoldi-/Lanczos-type method: $(I - C_j (\tilde{C}_j^T C_j)^{-1} \tilde{C}_j^T) S_j V_k = V_{k+1} \underline{T}_k$
- Biorthogonality condition: $\begin{bmatrix} \tilde{C}_j, W_k \end{bmatrix} \perp_b [C_j, V_k]$

Recycling subspace $C_j = S_j U_j$ and $\tilde{C}_j = S_j^{\mathrm{T}} \tilde{U}_j$

Recall:
$$S_j^T J = JS_j$$
 and $J = I$
• $J = J^T$: $\tilde{U}_j = JU_j$, $\tilde{C}_j = JC_j$
• $x_k = x_0 + (I - U_j \tilde{C}_j^T S_j) V_k y_k$



Multiple solution to linear system $S_j x_j = f_j$ with $f_j \in \mathbb{C}^{n \times p}$ and j = 1, ..., I.

Solution update:

$$x_k = x_0 + V_k y_k + U_j z$$

- Arnoldi-/Lanczos-type method: $(I - C_j (\tilde{C}_j^T C_j)^{-1} \tilde{C}_j^T) S_j V_k = V_{k+1} \underline{T}_k$
- Biorthogonality condition: $\begin{bmatrix} \tilde{C}_j, W_k \end{bmatrix} \perp_b [C_j, V_k]$

 $\begin{array}{l} \textbf{Recycling subspace} \\ C_{j} = S_{j}U_{j} \hspace{0.1 cm} \text{and} \hspace{0.1 cm} \tilde{C}_{j} = S_{j}^{\mathrm{T}}\tilde{U}_{j} \end{array}$

Recall:
$$S_j^{T}J = JS_j$$
 and $J = I$
• $J = J^{T}$: $\tilde{U}_j = JU_j$, $\tilde{C}_j = JC_j$
• $N_k = N_0 + (I - U_j \tilde{C}_j^{T} S_j) V_k y_k$



Multiple solution to linear system $S_j x_j = f_j$ with $f_j \in \mathbb{C}^{n \times p}$ and j = 1, ..., I.

Solution update:

$$x_k = x_0 + V_k y_k + U_j z$$

- Arnoldi-/Lanczos-type method: $(I - C_j (\tilde{C}_j^T C_j)^{-1} \tilde{C}_j^T) S_j V_k = V_{k+1} \underline{T}_k$
- Biorthogonality condition: $\begin{bmatrix} \tilde{C}_j, W_k \end{bmatrix} \perp_b [C_j, V_k]$

Recycling subspace $C_j = S_j U_j$ and $\tilde{C}_j = S_j^{\mathrm{T}} \tilde{U}_j$

Recall:
$$S_j^T J = JS_j$$
 and $J = I$
• $J = J^T$: $\tilde{U}_j = JU_j$, $\tilde{C}_j = JC_j$
• $x_k = x_0 + (I - U_j \tilde{C}_j^T S_j) V_k y_k$



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PCB circuit









- Frequency range: $[f_{min}, f_{max}] = [7.5, 10.0]$ GHz.
- Electric conductivity: $\sigma\equiv$ 0.
- Dimension of original model problem: *n* = 226458.



AORA-RK vs. AORA-MAX vs. AORA-H2



Figure: PCB circuit: $n = 226458 \rightarrow n_r = 20(80)$.



PCB circuit: mAORA method







(b)

Figure: PCB circuit: Modified adaptive-order rational Arnoldi method. $(n_r = 20)$



PCB circuit: mAORA method with rSQMR

Expansion point	# rSQMR	# SQMR
<i>s</i> ₁ = 5.44e+10ι	-	_
<i>s</i> ₂ = 6.28e+10ι	54, 24, 25, 24, 25, 24, 24, 24, 25, 25, 25	55
<i>s</i> ₃ = 6.05e+10ι	23, 20, 19, 20, 20, 20, 20, 20, 19	50
<i>s</i> ₄ = 5.84e+10ı	16, 13, 13, 13, 13, 13, 13	35
<i>s</i> ₅ = 4.71e+10ι	23, 21, 21, 18, 21	29
<i>s</i> ₆ = 5.73e+10ι	25, 11, 11, 10, 11	26

Table: PCB circuit: Iteration steps of rSQMR method in mAORA method.

Single preconditioning technique via

 $S = (s_*M_{\epsilon} + M_{\sigma}) + C(s_*M_{\mu})^{-1}C^{\mathrm{T}}$ with $s_* = \iota\sqrt{f_{\min}f_{\max}}$.

• Significantly smaller number of matrix-vector multiplications.



Outline

- Moment matching methods in model order reduction
- Efficient offline-stage of moment matching methods
 - Modified adaptive-order rational Arnoldi method
 - Recycling Krylov subspace methods
- Numerical experiments
- Conclusion



Conclusion



- Structure preservation.
- AORA-RK method.
- Divergence preservation.

- Modified rational Arnoldi method.
- Direct solver: ATLM.
- rSQMR method.



Conclusion



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