

Scientific Computing 1 10th Homework

Handout: 01/04/2017

Return: 01/13/2017

Exercise 1:

(3 Points)

Let $V, W \in \mathbb{R}^{n \times k}$, $k < n$, be two matrices of rank k . Show that

$$A = VW^T \in \mathbb{R}^{n \times n}$$

has exactly rank k .

Exercise 2:

(4 Points)

The LU decomposition is the main technique to solve moderate size linear systems and has an asymptotic flop count of $\frac{2}{3}n^3$ flops. Beside the classical scalar implementation nowadays block implementations with a blocking parameter r are used to exploit modern computer architectures.

- Show that the asymptotic flop count for the block LU decomposition is invariant under the chosen blocking parameter r .
- Compare the exact flop count (including the lower order terms) of the unblocked and the block implementation.

Exercise 3:

(14 Points)

The outer product Gaussian elimination defines one possible way to compute an LU decomposition of a matrix $A \in \mathbb{R}^{m \times n}$.

- Implement this algorithm **without** blocking techniques as a C function with the following signature:

```
int LU2(int m, int n, double * A, int lda);
```

The input `A` with leading dimension `lda` should be overwritten by its LU decomposition as shown in the lecture. The integer return value should be 0 if the matrix was decomposed successfully or a non zero value otherwise.

Hint: The appearing rank-1 updates can be performed using your `gemm` implementation from the previous exercises.

- Implement this algorithm **with** blocking techniques as a C function with the following signature:

```
int LU3(int m, int n, double * A, int lda, int r);
```

The input `A` with leading dimension `lda` should be overwritten by its LU decomposition as shown in the lecture. The blocking parameter `r` should be freely selectable by the caller. The integer return value should be 0 if the matrix was decomposed successfully or a non zero value otherwise. The the LU decomposition function from a.) to decompose the small panels.

Hint: The appearing rank-r updates can be performed using your `gemm` implementation from the previous exercises.

- c.) Write a solver function that takes a square LU decomposed matrix from a.) or b.), $L \in \mathbb{R}^{n \times n}$ and $U \in \mathbb{R}^{n \times n}$, and a right hand side $b \in \mathbb{R}^n$ as inputs and overwrites b with the solution of $LUx = b$. Use the BLAS function `DTRSV` or `DTRSM` for this purpose. The function header should be

```
void LU_solve(int n, double *LU, int ldlu, double *b);
```

where `ldlu` is the leading dimension of the matrix containing the LU decomposition.

- d.) Solve the following linear system to check your code:

$$\begin{pmatrix} 1 & \frac{1}{2} & \frac{1}{3} & \frac{1}{4} & \frac{1}{5} & \frac{1}{6} \\ \frac{1}{2} & \frac{1}{3} & \frac{1}{4} & \frac{1}{5} & \frac{1}{6} & \frac{1}{7} \\ \frac{1}{3} & \frac{1}{4} & \frac{1}{5} & \frac{1}{6} & \frac{1}{7} & \frac{1}{8} \\ \frac{1}{4} & \frac{1}{5} & \frac{1}{6} & \frac{1}{7} & \frac{1}{8} & \frac{1}{9} \\ \frac{1}{5} & \frac{1}{6} & \frac{1}{7} & \frac{1}{8} & \frac{1}{9} & \frac{1}{10} \\ \frac{1}{6} & \frac{1}{7} & \frac{1}{8} & \frac{1}{9} & \frac{1}{10} & \frac{1}{11} \end{pmatrix} x = \begin{pmatrix} \frac{49}{20} \\ \frac{223}{140} \\ \frac{341}{280} \\ \frac{2509}{2520} \\ \frac{2131}{2520} \\ \frac{20417}{27720} \end{pmatrix}$$

What can you recognize? Is the solution sufficiently accurate when you think about the possibilities of double precision floating point numbers?

Hint: A skeleton code with the data for d.) is provided on the lecture's website.

Exercise 4: (4 Points)

The Jacobi method is a classical splitting technique to solve linear systems iteratively. It splits a matrix $A \in \mathbb{R}^{n \times n}$ into the diagonal $D = \text{diag}\{a_{11}, a_{22}, \dots, a_{nn}\}$ and off-diagonal elements and repeatedly solves the equation

$$x_{i+1} = M_J x_i + D^{-1}b$$

where the iteration matrix is

$$M_J = D^{-1}(D - A).$$

Show that the Jacobi method converges to the solution $x = A^{-1}b$ if the matrix A is strictly diagonal-dominant, i.e.,

$$\sum_{j=1, j \neq i}^n |a_{ij}| < |a_{ii}|, \quad \forall i = 1, \dots, n.$$

Hint: Use Theorem 6.25.

Overall Points: 25