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# Block-diagonal preconditioning for optimal control problems constrained by PDEs with uncertain inputs 




#### Abstract

This paper is aimed at the efficient numerical simulation of optimization problems governed by either steady-state or unsteady partial differential equations involving random coefficients. This class of problems often leads to prohibitively high dimensional saddle point systems with tensor product structure, especially when discretized with the stochastic Galerkin finite element method. Here, we derive and analyze robust Schur complement-based block-diagonal preconditioners for solving the resulting stochastic optimality systems with all-at-once low-rank solvers. Moreover, we illustrate the effectiveness of our solvers with numerical experiments.


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Optimization problems constrained by deterministic steady-state partial differential equations (PDEs) are computationally challenging. This is even more so if the constraints are deterministic unsteady PDEs since one would then need to solve a system of PDEs coupled globally in time and space, and time-stepping methods quickly reach their limitations due to the enormous demand for storage [23, 28]. Yet, more challenging than the afore-mentioned are problems constrained by unsteady PDEs involving (countably many) parametric or uncertain inputs. This class of problems often leads to prohibitively high dimensional linear systems with Kronecker product structure, especially when discretized with the stochastic Galerkin finite element method (SGFEM). Moreover, a typical model for an optimal control problem with stochastic inputs (SOCP) will usually be used for the quantification of the statistics of the system response; this is a task that could in turn result in additional enormous computational expense.

Stochastic finite element-based solvers for a large range of PDEs with random data have been studied extensively [1, 2, 14, 25, 27, 32]. However, optimization problems constrained by PDEs with random inputs have, in our opinion, not yet received adequate attention. Some of the papers on these problems include [3, 13, 14, 15, 27, 31, While [13, 14] use SGFEM to study the existence and the uniqueness of control problems constrained by elliptic PDEs with random inputs, the emphasis in [3, 15, 31] is on solvers based on stochastic collocation methods for optimal control problems with random coefficients. Rosseel and Wells in 27 apply a one-shot method with both SGFEM and collocation approaches to an optimal control problem constrained by stochastic elliptic PDEs. One of their findings is that SGFEM generally exhibits superior performance compared to the stochastic collocation method, in the sense that, unlike SGFEM, the non-intrusivity property of the stochastic collocation method is lost when moments of the state variable appear in the cost functional, or when the control function is a deterministic function.
The fast convergence and other nice properties exhibited by SGFEM notwithstanding, the resulting large tensor-product algebraic systems associated with this intrusive approach unfortunately limits its attractiveness. Thus, for it to compete favourably with the sampling-based approaches, there is the need to develop efficient solvers for the resulting large linear systems. This is indeed the motivation for this work. More precisely, we apply an all-at-once approach, together with SGFEM, to two prototypical models, namely, optimization problems constrained by ( $a$ ) stationary diffusion equations, (b) unsteady diffusion equations, and in each of the two cases, both the constraint equations and the objective functional have uncertain inputs. As these problems pose increased computational complexity due to enormous memory requirements, we here focus specifically on efficient low-rank preconditioned iterative solvers for the resulting linear systems representing the Karush-Kuhn-Tucker (KKT) conditions. In particular, inspired by a state-of-the-art preconditioning strategy employed in the deterministic framework [24, 23, 28], we derive and analyze robust Schur complement-based blockdiagonal preconditioners which we use in conjunction with low-rank solvers for the efficient solution of the optimality systems.
In order to numerically simulate the above SOCP, we assume that the state, the control and the target (or the desired state) are analytic functions depending on the
uncertain parameters. This allows for a simultaneous generalized polynomial chaos approximation of these random functions in the SGFEM discretization of the models. However, we note here that, as pointed out in [27, problems in which the control is modeled as an unknown stochastic function constitute inverse problems and they are different from those with deterministic controls [35]. In the former, the stochastic properties of the control are unknown but will be computed. So, in most cases (as we assume in this work), the mean of the computed stochastic control could be considered as optimal. Depending on the application, the mean may not, in general, be the sought optimal control, though. Besides, computing the uncertainty in the system response might require additional computational challenges.

This paper is structured as follows. In Section 1, we present our problem statement and give an overview of the SGFEM on which we shall rely in the sequel. Section 2 discusses efficient solution of our first model problem, namely, an optimization problem governed by a steady-state diffusion equation with uncertain inputs. As an extension of the concepts discussed in Section 2, we proceed to Section 3 to introduce and analyze our preconditioning strategy for the unsteady analogue of the steady-state model. Furthermore, we here briefly review the tensor-train (TT) toolbox - a software which we shall use, in conjunction with MINRES, to solve our unsteady problems. Finally, Section 4 presents some numerical experiments to demonstrate the performance of our solvers.

## 1 Problem statement

In this paper, we study the numerical simulation of optimal control problems constrained by PDEs with uncertain coefficients. More precisely, we formulate our model problems as

$$
\begin{equation*}
\min _{y, u} \mathcal{J}(y, u) \text { subject to } c(y, u)=0 \tag{1}
\end{equation*}
$$

where the constraint equation $c(y, u)=0$ represents a PDE to be specified in the sequel, and

$$
\begin{equation*}
\mathcal{J}(y, u):=\frac{1}{2}\|y-\bar{y}\|_{L^{2}(\mathcal{D}) \otimes L^{2}(\Omega)}^{2}+\frac{\alpha}{2}\|\operatorname{std}(y)\|_{L^{2}(\mathcal{D})}^{2}+\frac{\beta}{2}\|u\|_{L^{2}(\mathcal{D}) \otimes L^{2}(\Omega)}^{2} \tag{2}
\end{equation*}
$$

is a cost functional of tracking-type. The functions $y, u$ and $\bar{y}$ are, in general, realvalued random fields representing, respectively, the state, the control and the prescribed target system response. We note here that $\bar{y}$ and $u$ could also be modeled deterministically. The positive constant $\beta$ in (2) represents the parameter for the penalization of the action of the control $u$, whereas $\alpha$ penalizes the standard deviation $\operatorname{std}(y)$ of the state $y$. The objective functional $\mathcal{J}(y, u)$ is a deterministic quantity with uncertain terms. In what follows, we shall focus mainly on distributed control problems, although we do believe that our discussion generalizes to boundary control problems as well.

Next, we recall that by a random field $z: \mathcal{D} \times \Omega \rightarrow \mathbb{R}$, we mean that $z(\mathbf{x}, \cdot)$ is a random variable defined on the complete probability space $(\Omega, \mathcal{F}, \mathbb{P})$ for each $\mathbf{x} \in \mathcal{D}$.

Moreover, for any random variable $g$ defined on $(\Omega, \mathcal{F}, \mathbb{P})$, the standard deviation $\operatorname{std}(g)$ of $g$ is given by

$$
\begin{equation*}
\operatorname{std}(g)=\left[\int_{\Omega}(g-\mathbb{E}(g))^{2} d \mathbb{P}(\omega)\right]^{\frac{1}{2}} \tag{3}
\end{equation*}
$$

where $\mathbb{E}(g)$ is the mean of $g$ defined by

$$
\begin{equation*}
\langle g\rangle:=\mathbb{E}(g)=\int_{\Omega} g d \mathbb{P}(\omega)<\infty . \tag{4}
\end{equation*}
$$

The Kronecker product Hilbert space $L^{2}(\mathcal{D}) \otimes L^{2}(\Omega)$ is endowed with the norm

$$
\|v\|_{L^{2}(\mathcal{D}) \otimes L^{2}(\Omega)}:=\left(\int_{\Omega}\|v(\cdot, \omega)\|_{L^{2}(\mathcal{D})}^{2} d \mathbb{P}(\omega)\right)^{\frac{1}{2}}<\infty
$$

where $L^{2}(\Omega):=L^{2}(\Omega, \mathcal{F}, \mathbb{P})$.

### 1.1 Representation of random inputs

In the spirit of [27], we consider two ways of representing the random fields which we shall use in the rest of the paper. In doing so, we will employ the so-called finite noise assumption, which states that a random field can be approximated with a prescribed finite number of random variables $\xi:=\left\{\xi_{1}, \xi_{2}, \ldots, \xi_{N}\right\}$, where $N \in \mathbb{N}$ and $\xi_{i}: \Omega \rightarrow \Gamma_{i} \subseteq \mathbb{R}$. We also make the simplifying assumption that each random variable is independent and characterized by a probability density function $\rho_{i}: \Gamma_{i} \rightarrow[0,1]$. Moreover, the random vector $\xi$ has a bounded joint probability density function $\rho: \Gamma \rightarrow \mathbb{R}^{+}$, where $\Gamma:=\prod_{i=1}^{N} \Gamma_{i} \subset \mathbb{R}^{N}$ and $\rho=\prod_{i=1}^{N} \rho_{i}\left(\xi_{i}\right)$. With these assumptions, DoobDynkin's Lemma, cf. [1], guarantees that one can therefore parametrically represent a random field $z(\mathbf{x}, \omega)$ in terms of the random vector $\xi$ instead of the random outcomes $\omega$. That is, we can now write $z\left(\mathbf{x}, \xi_{1}, \xi_{2}, \ldots, \xi_{N}\right)$. Besides, denoting the space of squareintegrable random variables with respect to the density $\rho$ by $L_{\rho}^{2}(\Gamma)$, we introduce the space $L^{2}(\mathcal{D}) \otimes L_{\rho}^{2}(\Gamma)$, equipped with the norm

$$
\begin{equation*}
\|v\|_{L^{2}(\mathcal{D}) \otimes L_{\rho}^{2}(\Gamma)}:=\left(\int_{\Gamma}\|v(\cdot, \xi)\|_{L^{2}(\mathcal{D})}^{2} \rho(\xi) d \xi\right)^{\frac{1}{2}}<\infty \tag{5}
\end{equation*}
$$

Similarly, using equations (3) and (4) we have

$$
\begin{equation*}
\operatorname{std}(g)=\left[\int_{\Gamma}(g(\xi)-\mathbb{E}(g(\xi)))^{2} \rho(\xi) d \xi\right]^{\frac{1}{2}} \quad \text { and } \quad\langle g\rangle=\int_{\Gamma} g(\xi) \rho(\xi) d \xi<\infty \tag{6}
\end{equation*}
$$

Furthermore, our cost functional $\mathcal{J}(y, u)$ now reads

$$
\begin{equation*}
\mathcal{J}(y, u):=\frac{1}{2}\|y-\bar{y}\|_{L^{2}(\mathcal{D}) \otimes L_{\rho}^{2}(\Gamma)}^{2}+\frac{\alpha}{2}\|\operatorname{std}(y)\|_{L^{2}(\mathcal{D})}^{2}+\frac{\beta}{2}\|u\|_{L^{2}(\mathcal{D}) \otimes L_{\rho}^{2}(\Gamma)}^{2} \tag{7}
\end{equation*}
$$

Suppose now that $z: \mathcal{D} \times \Omega \rightarrow \mathbb{R}$ is a random field with known continuous covariance function $C_{z}(\mathbf{x}, \mathbf{y})$. Then, one way to represent $z$ with a finite number of random variables is through a truncated Karhunen-Lòeve expansion (KLE):

$$
\begin{equation*}
z_{N}(\mathbf{x}, \xi(\omega))=\mathbb{E}[z](\mathbf{x})+\sigma_{z} \sum_{i=1}^{N} \sqrt{\lambda_{i}} \varphi_{i}(\mathbf{x}) \xi_{i}(\omega) \tag{8}
\end{equation*}
$$

where $\sigma_{z}$ is the standard deviation of $z$, the random variables $\left\{\xi_{i}\right\}_{i=1}^{N}$ are centered, normalized and uncorrelated with

$$
\xi_{i}(\omega)=\frac{1}{\sigma_{z} \sqrt{\lambda_{i}}} \int_{\mathcal{D}}(z(\mathbf{x}, \xi(\omega))-\mathbb{E}[z](\mathbf{x})) \varphi_{i}(\mathbf{x}) d \mathbf{x}
$$

and $\left\{\lambda_{i}, \varphi_{i}\right\}$ is the set of eigenvalues and eigenfunctions corresponding to $C_{z}(\mathbf{x}, \mathbf{y})$, that is,

$$
\int_{\mathcal{D}} C_{z}(\mathbf{x}, \mathbf{y}) \varphi_{i}(\mathbf{y}) d \mathbf{y}=\lambda_{i} \varphi_{i}(\mathbf{x})
$$

The eigenfunctions $\left\{\varphi_{i}\right\}$ form a complete orthogonal basis in $L^{2}(\mathcal{D})$. The eigenvalues $\left\{\lambda_{i}\right\}$ form a sequence of non-negative real numbers decreasing to zero. The series (8) represents the best $N$-term approximation of $z$ and converges in $L^{2}(\mathcal{D}) \otimes L_{\rho}^{2}(\Gamma)$, due to

$$
\sum_{i=1}^{\infty} \lambda_{i}=\int_{\Gamma} \int_{\mathcal{D}} \rho(z(\mathbf{x}, \xi(\omega))-\mathbb{E}[z](\mathbf{x}))^{2} d \mathbf{x} d \xi
$$

Alternatively, one can approximate the random field of interest $z$ using a truncated generalized polynomial chaos expansion (PCE):

$$
\begin{equation*}
z_{P}(\mathbf{x}, \omega)=\sum_{j=0}^{P-1} z_{j}(\mathbf{x}) \psi_{j}(\xi(\omega)) \tag{9}
\end{equation*}
$$

where $z_{j}$, the deterministic modes of the expansion, are given by

$$
z_{j}(\mathbf{x}, \omega)=\frac{\left\langle z(\mathbf{x}, \xi(\omega)) \psi_{j}(\xi)\right\rangle}{\left\langle\psi_{j}^{2}(\xi)\right\rangle}
$$

$\left\{\psi_{j}\right\}_{j=0}^{P-1}$ are $N$-variate orthogonal polynomials of order $n$ and

$$
\begin{equation*}
P=1+\sum_{k=1}^{n} \frac{1}{k!} \prod_{j=0}^{k-1}(N+j)=\frac{(N+n)!}{N!n!} . \tag{10}
\end{equation*}
$$

The $N$-variate orthogonal polynomials $\left\{\psi_{j}\right\}_{j=0}^{P-1}$ satisfy

$$
\begin{equation*}
\left\langle\psi_{0}(\xi)\right\rangle=1, \quad\left\langle\psi_{j}(\xi)\right\rangle=0, j>0, \quad\left\langle\psi_{j}(\xi) \psi_{k}(\xi)\right\rangle=\left\langle\psi_{j}^{2}(\xi)\right\rangle \delta_{j k}, \tag{11}
\end{equation*}
$$

with

$$
\begin{equation*}
\left\langle\psi_{j}(\xi)\right\rangle=\int_{\omega \in \Omega} \psi_{j}(\xi(\omega)) d \mathbb{P}(\omega)=\int_{\xi \in \Gamma} \psi_{j}(\xi) \rho(\xi) d \xi \tag{12}
\end{equation*}
$$

Moreover, it can be easily shown then that

$$
\begin{equation*}
\mathbb{E}\left[z_{P}\right](\mathbf{x})=z_{0}(\mathbf{x}) \text { and } \operatorname{Var}\left[z_{P}\right](\mathbf{x})=\sum_{i=1}^{P-1} z_{i}^{2}(\mathbf{x})\left\langle\psi_{i}^{2}(\xi)\right\rangle \tag{13}
\end{equation*}
$$

Observe then that since $\operatorname{Var}(y)=[\operatorname{std}(y)]^{2}$, it follows immediately from 13) that

$$
\begin{equation*}
\|\operatorname{std}(y)\|_{L^{2}(\mathcal{D})}^{2}=\int_{\mathcal{D}} \mathbb{V} \operatorname{ar}[y](\mathbf{x}) d \mathbf{x}=\int_{\mathcal{D}} \sum_{i=1}^{P-1} y_{i}^{2}(\mathbf{x})\left\langle\psi_{i}^{2}(\xi)\right\rangle d \mathbf{x} \tag{14}
\end{equation*}
$$

In this contribution, we shall rely on the SGFEM for the spatial and stochastic discretizations, see e.g. [1, 14, 25, 27, 32], and our exposition here follows closely the framework in [27]. To this end, we note that the KLE and PCE representations are quite essential in the SGFEM discretizations of the optimal control problems discussed in the sequel [14, 27. In a nutshell, we recall that the SGFEM is a spectral approach in which one seeks $y$ and $u$ in a finite-dimensional subspace of the Hilbert space $H_{0}^{1}(\mathcal{D}) \otimes L_{\rho}^{2}(\Gamma)$, consisting of tensor products of deterministic functions defined on the spatial domain and stochastic functions defined on the probability space. More precisely, suppose first that $V_{h} \subset H_{0}^{1}(\mathcal{D})$ is a space of standard Lagrangian finite element functions on a partition $\mathcal{T}$ into triangles (or rectangles) of the domain $\mathcal{D}$ defined by

$$
V_{h}:=\left\{v_{h} \in H_{0}^{1}(\mathcal{D}): v_{h} \in P_{k}(\Xi) \forall \Xi \in \mathcal{T}\right\},
$$

where $\Xi \in \mathcal{T}$ is a cell and $P_{k}$ is the space of Lagrangian polynomials of degree $k$. In particular, let $V_{h}=\operatorname{span}\left\{\phi_{j}(\mathbf{x}), j=1, \ldots, J\right\}$. Moreover, let $Y_{n} \subset L_{\rho}^{2}(\Gamma)$ be such that $Y_{n}:=\operatorname{span}\left\{\psi_{i}(\xi): i \in \mathcal{I}\right\}$, where $\mathcal{I}:=\left\{i=\left(i_{1}, \ldots, i_{N}\right) \in \mathbb{N}^{N}:|i| \leq n\right\}$. That is, $Y_{n}$ is a set of all $N$-variate orthogonal polynomials of degree at most $n$, whereas $\mathcal{I}$ is a set of all multi-indices of length $N$ satisfying $|i| \leq n$. It can then be shown that

$$
\operatorname{dim}\left(Y_{n}\right)=\operatorname{dim}(\mathcal{I})=\binom{N+n}{N}
$$

which is precisely $P$ given in (10). Hence, it turns out that there exists a bijection $\mu:\{1, \ldots, P\} \rightarrow \mathcal{I}$ that assigns a unique integer $i$ to each multi-index $\mu(i) \in \mathcal{I}$.
To illustrate here how the space $Y_{n}$ is constructed [25], consider the case of uniform random variables with $N=2$ and $n=3$. Then $Y_{n}$ is a set of two-dimensional Legendre polynomials (products of a univariate Legendre polynomial in $\xi_{1}$ and a univariate Legendre polynomial in $\xi_{2}$ ) of degree less than or equal to three. Each of the basis functions is associated with a multi-index $\nu=\left(\nu_{1}, \nu_{2}\right)$, where the components represent the degrees of the polynomials in $\xi_{1}$ and $\xi_{2}$. Since the total degree of the polynomial is three, we have the possibilities $\nu=(0,0),(1,0),(2,0),(3,0),(0,1),(1,1),(2,1),(0,2)$,
$(1,2)$, and $(0,3)$. Since the univariate Legendre polynomials of degrees $0,1,2,3$ are $L_{0}(x)=1, L_{1}(x)=x, L_{2}(x)=\frac{1}{2}\left(3 x^{2}-1\right)$, and $L_{3}(x)=\frac{1}{2}\left(5 x^{3}-3 x\right)$, we have

$$
\begin{aligned}
Y_{n}= & \operatorname{span}\left\{\psi_{i}(\xi)\right\}_{i=0}^{9} \\
= & \left\{1, \xi_{1}, \frac{1}{2}\left(3 \xi_{1}^{2}-1\right), \frac{1}{2}\left(5 \xi_{1}^{3}-3 \xi_{1}\right), \xi_{2}, \xi_{1} \xi_{2}, \frac{1}{2}\left(3 \xi_{1}^{2}-1\right) \xi_{2}, \frac{1}{2}\left(3 \xi_{2}^{2}-1\right),\right. \\
& \left.\frac{1}{2} \xi_{1}\left(3 \xi_{2}^{2}-1\right), \frac{1}{2}\left(5 \xi_{2}^{3}-3 \xi_{2}\right)\right\} .
\end{aligned}
$$

So, the SGFEM essentially entails performing a Galerkin projection onto $W_{h n}:=$ $V_{h} \otimes Y_{n} \subset H_{0}^{1}(\mathcal{D}) \otimes L_{\rho}^{2}(\Gamma)$ using basis functions $r_{h n}$ of the form

$$
\begin{equation*}
r_{h n}=\sum_{j=1}^{J} \sum_{i \in \mathcal{I}} r_{i j} \phi_{j}(\mathbf{x}) \psi_{i}(\xi), \tag{15}
\end{equation*}
$$

where $r_{i j}$ is a degree of freedom. In this paper, $\left\{\phi_{j}\right\}$ are $\mathbf{Q}_{1}$ finite elements, whereas $\left\{\psi_{i}\right\}$ are multi-dimensional Legendre polynomials.
We now proceed to Section 2 to present our first SOCP whose constraint is a stationary diffusion equation. The idea is to use this model to motivate our discussion on the solvers for a time-dependent model problem in Section 3 .

## 2 A control problem with stationary diffusion equation

Our first SOCP consists now in minimizing the cost functional $\mathcal{J}(y(\mathbf{x}, \omega), u(\mathbf{x}, \omega))$ defined in (2) such that, $\mathbb{P}$-almost surely, the following linear elliptic diffusion equation holds

$$
\left\{\begin{align*}
-\nabla \cdot(a(\mathbf{x}, \omega) \nabla y(\mathbf{x}, \omega)) & =u(\mathbf{x}, \omega), \text { in } \mathcal{D} \times \Omega,  \tag{16}\\
y(\mathbf{x}, \omega) & =0, \text { on } \partial \mathcal{D} \times \Omega
\end{align*}\right.
$$

where $a: \mathcal{D} \times \Omega \rightarrow \mathbb{R}$ is a random coefficient field and the forcing term on the right hand side $u: \mathcal{D} \times \Omega \rightarrow \mathbb{R}$ denotes a random control function. Furthermore, we assume that

$$
\begin{equation*}
u \in L^{2}(\mathcal{D}) \otimes L^{2}(\Omega) \tag{17}
\end{equation*}
$$

and that there exist positive constants $a_{\min }$ and $a_{\max }$ such that

$$
\begin{equation*}
\mathbb{P}\left(\omega \in \Omega: a(\mathbf{x}, \omega) \in\left[a_{\min }, a_{\max }\right], \forall \mathbf{x} \in \mathcal{D}\right)=1 \tag{18}
\end{equation*}
$$

For the weak formulation of the forward problem 16 , we seek $y \in H_{0}^{1}(\mathcal{D}) \otimes L^{2}(\Omega)$ such that, $\mathbb{P}$-almost surely,

$$
\begin{equation*}
\mathcal{B}(y, v)=\ell(u, v), \quad \forall v \in H_{0}^{1}(\mathcal{D}) \otimes L^{2}(\Omega) \tag{19}
\end{equation*}
$$

where the bilinear form $\mathcal{B}(\cdot, \cdot)$ is given by

$$
\mathcal{B}(y, v):=\int_{\Omega} \int_{\mathcal{D}} a(\mathbf{x}, \omega) \nabla_{\mathbf{x}} y(\mathbf{x}, \omega) \cdot \nabla_{\mathbf{x}} v(\mathbf{x}, \omega) d \mathbf{x} d \mathbb{P}(\omega), \quad v, y \in H_{0}^{1}(\mathcal{D}) \otimes L^{2}(\Omega)
$$

and

$$
\ell(u, v):=\int_{\Omega} \int_{\mathcal{D}} u(\mathbf{x}, \omega) v(\mathbf{x}, \omega) d \mathbf{x} d \mathbb{P}(\omega), \quad v, u \in H_{0}^{1}(\mathcal{D}) \otimes L^{2}(\Omega)
$$

The following existence and uniqueness result of the solution $y$ to proved in, for instance, 14 follows from the Lax-Milgram Lemma [5].
Theorem 1. Under the assumptions (17) and (18), there exists a unique solution $y \in H_{0}^{1}(\mathcal{D}) \otimes L^{2}(\Omega)$ such that, $\mathbb{P}$-almost surely, 19$)$ holds.

Recasting the above SOCP given by $\sqrt{2}$ and 16 into a saddle-point formulation, Chen and Quarteroni in [6] prove the existence and uniqueness of its solution. More precisely, the following result holds.
Theorem 2. [6, Theorem 3.5] Let (17) and (18) be satisfied and let $\alpha=0$ in (7). Then, there exists a unique solution $(y, u)$ to the SOCP (7) and (16) satisfying the stochastic optimality conditions

$$
\begin{aligned}
\mathcal{B}(y, v) & =\ell(u, v), & & v \in H_{0}^{1}(\mathcal{D}) \otimes L^{2}(\Omega), \\
\ell(\beta u-f, w) & =0, & & w \in L^{2}(\mathcal{D}) \otimes L^{2}(\Omega), \\
\mathcal{B}^{\prime}(y, r)+\ell(y, r) & =\ell(\bar{y}, r), & & r \in H_{0}^{1}(\mathcal{D}) \otimes L^{2}(\Omega),
\end{aligned}
$$

where $f$ is the adjoint variable or Lagrangian parameter associated with the optimal solution $(y, u)$, and $\mathcal{B}^{\prime}$ is the adjoint bilinear form of $\mathcal{B}$; that is, $\mathcal{B}^{\prime}(y, r)=\mathcal{B}(r, y)$.

We note here that the cost functional considered in [6, 14] does not include $\|\operatorname{std}(y)\|_{L^{2}(\mathcal{D})}^{2}$. But then, their results extend to the more general form of $\mathcal{J}(y, u)$ considered in this paper due to the Frechét differentiability of $\|\operatorname{std}(y)\|_{L^{2}(\mathcal{D})}^{2}$; see, for example, [27].

As our major concern in this paper is to study efficient solvers resulting from the discretization of our model problems, we proceed next to recall the two common approaches in the literature to solve these optimization problems [29, 30. The first method is the so-called optimize-then-discretize (OTD) approach. Here, one essentially considers the infinite-dimensional problem, writes down the first order conditions and then discretizes the first order conditions. An alternative strategy, namely, the discretize-then-optimize (DTO) approach involves discretizing the problem first and then building a discrete Lagrangian functional with the corresponding first order conditions. The commutativity of DTO and OTD methods when applied to optimal control problems constrained by PDEs has been a subject of debate in recent times (see [17] for an overview). In what follows, we will adopt the DTO strategy because, for the SOCPs considered in this paper, it leads to a symmetric saddle point linear system which fits in nicely with our preconditioning strategy.
To discretize the SOCP given by (2) and (16) using the SGFEM, consider first the constraint 16. Given a basis for $W_{h n}:=V_{h} \otimes Y_{n} \subset H_{0}^{1}(\mathcal{D}) \otimes L_{\rho}^{2}(\Gamma)$ and a truncated

KLE representation $a_{N}(\mathbf{x}, \xi)$ (cf. (8)) of the random field $a$ satisfying (18), we now seek a finite dimensional $y_{h n}, u_{h n} \in W_{h n}$, satisfying

$$
\begin{equation*}
\int_{\Gamma} \int_{\mathcal{D}} a_{N}(\mathbf{x}, \xi) \nabla y_{h n} \cdot \nabla v \rho(\xi) d \mathbf{x} d \xi=\int_{\Gamma} \int_{\mathcal{D}} u_{h n} v \rho(\xi) d \mathbf{x} d \xi \tag{20}
\end{equation*}
$$

$\forall v \in W_{h n}$. Expanding $y_{h n}, u_{h n}$ and the test functions in the chosen basis in 20, we see that

$$
y_{h n}=\sum_{k=0}^{P-1} \sum_{j=1}^{J} y_{j k} \phi_{j}(\mathbf{x}) \psi_{k}(\xi)=\sum_{k=0}^{P-1} y_{k} \psi_{k}(\xi)
$$

and

$$
u_{h n}=\sum_{k=0}^{P-1} \sum_{j=1}^{J} u_{j k} \phi_{j}(\mathbf{x}) \psi_{k}(\xi)=\sum_{k=0}^{P-1} u_{k} \psi_{k}(\xi),
$$

yield the following linear system of dimension $J P \times J P$

$$
\begin{equation*}
\mathcal{K} \mathbf{y}=\mathcal{M} \mathbf{u} \tag{21}
\end{equation*}
$$

with block structure, where the blocks $\mathcal{K}_{p, q}$ of the stochastic Galerkin matrix $\mathcal{K}$ are linear combinations of $N+1$ weighted stiffness matrices of dimension $J$, with each of them having the same sparsity pattern equivalent to that of the corresponding deterministic problem. More specifically, for $p, q=0, \ldots, P-1$, we have

$$
\begin{equation*}
\mathcal{K}_{p, q}=\left\langle\psi_{p}(\xi) \psi_{q}(\xi)\right\rangle K_{0}+\sum_{i=1}^{N}\left\langle\xi_{i} \psi_{p}(\xi) \psi_{q}(\xi)\right\rangle K_{i}, \tag{22}
\end{equation*}
$$

and

$$
\mathcal{M}_{p, q}=\left\langle\psi_{p}(\xi) \psi_{q}(\xi)\right\rangle M,
$$

where the mass matrix $M \in \mathbb{R}^{J \times J}$ and the stiffness matrices $K_{i} \in \mathbb{R}^{J \times J}, i=$ $0,1, \ldots, N$, are given, respectively, by

$$
\begin{align*}
M(j, k) & =\int_{\mathcal{D}} \phi_{j}(\mathbf{x}) \phi_{k}(\mathbf{x}) d \mathbf{x}  \tag{23}\\
K_{0}(j, k) & =\int_{\mathcal{D}} \mathbb{E}[a](\mathbf{x}) \nabla \phi_{j}(\mathbf{x}) \nabla \phi_{k}(\mathbf{x}) d \mathbf{x}  \tag{24}\\
K_{i}(j, k) & =\sigma_{a} \sqrt{\lambda_{i}} \int_{\mathcal{D}} \varphi_{i}(\mathbf{x}) \nabla \phi_{j}(\mathbf{x}) \nabla \phi_{k}(\mathbf{x}) d \mathbf{x} \tag{25}
\end{align*}
$$

where we assume that $\mathbb{E}[a]>0$, so that $K_{0}$ is symmetric and positive definite. The block $K_{0}$ captures the mean information in the model, whereas the other blocks $K_{i}, i=$ $1, \ldots, N$, represent fluctuations in the model. In Kronecker product notation, one obtains

$$
\begin{equation*}
\mathcal{K}:=G_{0} \otimes K_{0}+\sum_{i=0}^{N} G_{i} \otimes K_{i}, \quad \mathcal{M}:=G_{0} \otimes M \tag{26}
\end{equation*}
$$

where

$$
\left\{\begin{array}{l}
G_{0}=\operatorname{diag}\left(\left\langle\psi_{0}^{2}\right\rangle,\left\langle\psi_{1}^{2}\right\rangle, \ldots,\left\langle\psi_{P-1}^{2}\right\rangle\right)  \tag{27}\\
G_{i}(j, k)=\left\langle\xi_{i} \psi_{j} \psi_{k}\right\rangle, \quad i=1, \ldots, N
\end{array}\right.
$$

due to the orthogonality of the stochastic basis functions with respect to the probability measure of the distribution of the chosen random variables (cf. 11). Moreover, $\mathcal{K}$ is highly sparse as many of the sums in 22 are zero.

Similarly, applying SGFEM to the cost function (7), taking into account the expression (14), leads to

$$
\begin{equation*}
\frac{1}{2}(\mathbf{y}-\overline{\mathbf{y}})^{T} \mathcal{M}(\mathbf{y}-\overline{\mathbf{y}})+\frac{\alpha}{2} \mathbf{y}^{T} \mathcal{M}_{t} \mathbf{y}+\frac{\beta}{2} \mathbf{u}^{T} \mathcal{M} \mathbf{u} \tag{28}
\end{equation*}
$$

where

$$
\begin{equation*}
\mathcal{M}_{t}:=T \otimes M, \quad T:=\operatorname{diag}\left(0,\left\langle\psi_{1}^{2}\right\rangle, \ldots,\left\langle\psi_{P-1}^{2}\right\rangle\right) \tag{29}
\end{equation*}
$$

Our discrete SOCP now is to minimize (28) subject to 21. The Lagrangian functional $\mathcal{L}$ of this optimization problem is given by

$$
\mathcal{L}(\mathbf{y}, \mathbf{u}, \mathbf{f}):=\frac{1}{2}(\mathbf{y}-\overline{\mathbf{y}})^{T} \mathcal{M}(\mathbf{y}-\overline{\mathbf{y}})+\frac{\alpha}{2} \mathbf{y}^{T} \mathcal{M}_{t} \mathbf{y}+\frac{\beta}{2} \mathbf{u}^{T} \mathcal{M} \mathbf{u}+\mathbf{f}^{T}(-\mathcal{K} \mathbf{y}+\mathcal{M} \mathbf{u})
$$

where $\mathbf{f}$ denotes the Lagrangian multiplier or adjoint associated with the constraint. Now, applying the first order conditions to the Lagrangian yields the following optimality system

$$
\underbrace{\left[\begin{array}{ccc}
\mathcal{M}_{\alpha} & 0 & -\mathcal{K}^{T}  \tag{30}\\
0 & \beta \mathcal{M} & \mathcal{M}^{T} \\
-\mathcal{K} & \mathcal{M} & 0
\end{array}\right]}_{:=\mathcal{A}}\left[\begin{array}{l}
\mathbf{y} \\
\mathbf{u} \\
\mathbf{f}
\end{array}\right]=\left[\begin{array}{c}
\mathcal{M} \overline{\mathbf{y}} \\
\mathbf{0} \\
\mathbf{d}
\end{array}\right]
$$

where

$$
\begin{align*}
\mathcal{M}_{\alpha} & =\mathcal{M}+\alpha \mathcal{M}_{t} \\
& =\left(G_{0} \otimes M\right)+\alpha(T \otimes M) \\
& =G_{\alpha} \otimes M \tag{31}
\end{align*}
$$

with $G_{\alpha}:=G_{0}+\alpha T$, so that

$$
G_{\alpha}(j, k)= \begin{cases}\left\langle\psi_{0}^{2}\right\rangle, & \text { if } j=k=0  \tag{32}\\ (1+\alpha)\left\langle\psi_{j}^{2}\right\rangle, & \text { if } j=k=1,2, \ldots, P-1, \\ 0, & \text { otherwise }\end{cases}
$$

We note from (26), (31) and (32) that if $\alpha=0$, then $G_{\alpha}=G_{0}$ and, hence, $\mathcal{M}_{\alpha}=$ $\mathcal{M}$. Moreover, we assume that the parameter $N$ in the KLE of the random input $a$
is chosen such that $\mathcal{K}$ stays symmetric and positive definite [25]. The vector $\mathbf{d}:=$ $\operatorname{diag}\left(G_{0}\right) \otimes \tilde{\mathbf{d}}$, where $\tilde{\mathbf{d}}$ represents, in general, contributions from boundary conditions with respect to the spatial discretization. The system (30) is usually of huge dimension. As a result, the use of direct solvers for the system is out of the question. In what follows, we consider efficient iterative solvers instead. First, however, we discuss our preconditioning strategies.

### 2.1 Preconditioning the optimality system

Now, observe that the optimality system (30) is of saddle point form 9:

$$
\mathcal{A}=\left[\begin{array}{cc}
A & B^{T}  \tag{33}\\
B & 0
\end{array}\right]
$$

where

$$
A=\left[\begin{array}{cc}
\mathcal{M}_{\alpha} & 0  \tag{34}\\
0 & \beta \mathcal{M}
\end{array}\right], \quad B=\left[\begin{array}{ll}
-\mathcal{K} & \mathcal{M}
\end{array}\right]
$$

where $A$ is symmetric and positive definite and $B$ has full row rank. An appropriate Krylov subspace solver for the indefinite saddle point system is the MINRES algorithm originally proposed by Paige and Saunders in [22]. However, effectiveness of iterative solvers requires a suitable preconditioner. That is, we need a matrix $\mathcal{P}$ such that $\mathcal{P}^{-1} \mathcal{A}$ has better spectral properties (essentially, clustered eigenvalues) and for which $\mathcal{P}^{-1} \mathbf{v}$ is cheap to compute for any vector $\mathbf{v}$ of appropriate dimension. In what follows, we discuss our preconditioning strategies for solving (30).

Throughout this paper, we will focus mainly on block diagonal preconditioners; that is, preconditioners of the form

$$
\mathcal{P}=\left[\begin{array}{cc}
A & 0 \\
0 & S
\end{array}\right]
$$

where $S=B A^{-1} B^{T}$ is the (negative) Schur complement. It has been shown [18] that this choice of preconditioner indeed yields three distinct eigenvalues $\left\{\frac{1-\sqrt{5}}{2}, 1, \frac{1+\sqrt{5}}{2}\right\}$ of $\mathcal{P}^{-1} \mathcal{A}$. Hence, any Krylov subspace method with optimality property, such as MINRES, will terminate after at most three iterations. With this preconditioner, we specifically have

$$
\mathcal{P}=\left[\begin{array}{ccc}
\mathcal{M}_{\alpha} & 0 & 0  \tag{35}\\
0 & \beta \mathcal{M} & 0 \\
0 & 0 & S
\end{array}\right]
$$

where

$$
\begin{equation*}
S=\mathcal{K} \mathcal{M}_{\alpha}^{-1} \mathcal{K}+\frac{1}{\beta} \mathcal{M} \tag{36}
\end{equation*}
$$

since $\mathcal{K}$ and $\mathcal{M}$ are symmetric. We note here that (35) is only an ideal preconditioner for our saddle point system (30) in the sense that it is not cheap to solve the system with it. In practice, one often has to approximate the three diagonal blocks in order to use $\mathcal{P}$ with MINRES. An effective approach to approximate blocks $(1,1)$ and $(2,2)$ is the application of Chebyshev semi-iteration to the mass matrices in each of the two blocks 34. More specifically, for a given system involving a mass matrix $M \mathbf{x}=\mathbf{b}$, the Chebyshev semi-iteration, as given by Algorithm 11, is used to speed up a relaxed Jacobi iteration:

$$
\mathbf{x}_{k+1}=H \mathbf{x}_{k}+\mathbf{g}
$$

where $H=I-\theta D_{0}^{-1} M, \mathbf{g}=\theta D_{0}^{-1} \mathbf{b}, D_{0}=\operatorname{diag}(M)$. The optimal relaxation parameter $\theta$ must be chosen in such a way that the the spectrum of the matrix $H$ is symmetric about the origin. For instance, for a mesh of square $\mathbf{Q}_{1}$ elements in 2 dimensions, $\lambda\left(D_{0}^{-1} M\right) \subset[1 / 4,9 / 4] ;$ moreover, if $\theta=4 / 5$, then we get $\lambda(H) \subset[-4 / 5,4 / 5]$; see e.g. 33.

```
Algorithm 1 Chebyshev semi-iterative algorithm for \(\ell\) steps
    Set \(D_{0}=\operatorname{diag}(M)\).
    Set relaxation parameter \(\theta\).
    Compute \(\mathbf{g}=\theta D_{0}^{-1} \mathbf{b}\).
    Set \(H=I-\theta D_{0}^{-1} M\) (this can be used implicitly).
    Set \(\mathbf{x}_{0}=0\) and \(\mathbf{x}_{k}=H \mathbf{x}_{k-1}+g\).
    Set \(c_{0}=2\) and \(c_{1}=\theta\).
    for \(k=1, \ldots, l\) do
        \(c_{k+1}=\theta c_{k}-\frac{1}{4} c_{k-1}\).
        \(\vartheta_{k+1}=\theta \frac{c_{k}}{c_{k+1}}\).
        \(\mathbf{x}_{k+1}=\vartheta_{k+1}\left(H \mathbf{x}_{k}+\mathbf{g}-\mathbf{x}_{k-1}\right)+\mathbf{x}_{k-1}\).
    end for
```

Approximating the Schur complement $S$, that is, block $(3,3)$ poses more difficulty, however. One possibility [26] is to approximate $S$ by dropping the term $\frac{1}{\beta} \mathcal{M}$ to obtain

$$
\begin{equation*}
S_{0}:=\mathcal{K} \mathcal{M}_{\alpha}^{-1} \mathcal{K}^{T} \tag{37}
\end{equation*}
$$

An alternative and more robust approach, which we adopt here and in the rest of this paper, was proposed in [24] in the context of deterministic optimal control problems. In this case, $S$ is approximated by a matrix $S_{1}$ of the form

$$
\begin{equation*}
S_{1}=\left(\mathcal{K}+\mathcal{M}_{u}\right) \mathcal{M}_{\alpha}^{-1}\left(\mathcal{K}+\mathcal{M}_{u}\right)^{T} \tag{38}
\end{equation*}
$$

where $\mathcal{M}_{u}$ is determined by 'matching' the terms in the expressions for $S_{1}$ and $S$ as given, respectively, in (38) and (36). More precisely, we ignore the cross terms (that is, $\left.\mathcal{K} \mathcal{M}_{\alpha}^{-1} \mathcal{M}_{u}+\mathcal{M}_{u} \mathcal{M}_{\alpha}^{-1} \mathcal{K}\right)$ in the expansion of $S_{1}$ to get

$$
\begin{equation*}
\mathcal{M}_{u} \mathcal{M}_{\alpha}^{-1} \mathcal{M}_{u}=\frac{1}{\beta} \mathcal{M}=\frac{1}{\beta} \mathcal{M} \mathcal{M}^{-1} \mathcal{M} \tag{39}
\end{equation*}
$$

Now, observe from (27), (31) and (32), we have that $\mathcal{M}_{\alpha}=G_{\alpha} \otimes M$. Moreover, note that ideally in (7), we have $\alpha \geq 0$. So, to derive an approximation to $S_{1}$, we consider first of all the case $\alpha=0$. In this case, it is easy to see that holds if we set

$$
\begin{equation*}
\mathcal{M}_{u}=\frac{1}{\sqrt{\beta}} \mathcal{M} \tag{40}
\end{equation*}
$$

since $\mathcal{M}_{\alpha}=\mathcal{M}$. If $\alpha>0$, then we apply the following trick. We proceed first to replace in equation 32 the $(0,0)$ entry in the diagonal matrix $G_{\alpha}$ by $(1+\alpha)\left\langle\psi_{0}^{2}\right\rangle$, so that we can then obtain

$$
\mathcal{M}_{\alpha}=G_{\alpha} \otimes M \approx(1+\alpha) G_{0} \otimes M=(1+\alpha) \mathcal{M}
$$

It turns out then that holds if and only if

$$
\mathcal{M}_{u}=\sqrt{\frac{1+\alpha}{\beta}} \mathcal{M}
$$

with which we recover 40 for $\alpha=0$. Hence, we have

$$
\begin{equation*}
S_{1}=\left(\mathcal{K}+\sqrt{\frac{1+\alpha}{\beta}} \mathcal{M}\right) \mathcal{M}_{\alpha}^{-1}\left(\mathcal{K}+\sqrt{\frac{1+\alpha}{\beta}} \mathcal{M}\right)^{T} \tag{41}
\end{equation*}
$$

We point out here that the expression for $\mathcal{M}_{u}$ implies that the ignored cross terms are $\mathcal{O}\left(\beta^{-1 / 2}\right)$ instead of $\mathcal{O}\left(\beta^{-1}\right)$ in (37).

The effectiveness of the iterative solver used to solve our KKT system depends to a large extent on how well the approximation $S_{1}$ represents the exact Schur complement. To measure this, we need to consider the eigenvalues of the preconditioned Schur complement $S_{1}^{-1} S$. In what follows, we proceed to establish the spectrum $\lambda\left(S_{1}^{-1} S\right)$ of $S_{1}^{-1} S$ by examining the Raleigh quotient

$$
R(x):=\frac{x^{T} S x}{x^{T} S_{1} x},
$$

for any non-zero vector $x$ of appropriate dimension. We shall rely on the following result on positive definite matrices.

Proposition 1. [19, Theorem 2] Let $X=A B+B A$, where $A$ and $B$ are positive definite, Hermitian square matrices. Then, $X$ is positive definite if

$$
\kappa(B)<\left(\frac{\sqrt{\kappa(A)}+1}{\sqrt{\kappa(A)}-1}\right)^{2}
$$

where $\kappa(Y)$ represents the spectral condition number of the matrix $Y$.
We can now prove the main result of this section.

Theorem 3. Let $\alpha \in[0,+\infty)$. Then, the eigenvalues of $S_{1}^{-1} S$ satisfy

$$
\begin{equation*}
\lambda\left(S_{1}^{-1} S\right) \subset\left[\frac{1}{2(1+\alpha)}, 1\right), \quad \forall \alpha<\left(\frac{\sqrt{\kappa(\mathcal{C})}+1}{\sqrt{\kappa(\mathcal{C})}-1}\right)^{2}-1 \tag{42}
\end{equation*}
$$

where $\mathcal{C}=\mathcal{M}^{-1 / 2} \mathcal{K} \mathcal{M}^{-1 / 2}$.
Proof. Suppose that $\alpha \in[0,+\infty)$. Define the diagonal matrices $\Upsilon$ and $\mathcal{E}_{\alpha}$ by

$$
\begin{equation*}
\Upsilon=\operatorname{diag}\left(0, I_{P-1}\right) \text { and } \mathcal{E}_{\alpha}=\left(I_{P}+\alpha \Upsilon\right) \otimes I_{J} \tag{43}
\end{equation*}
$$

where $I_{n}$ denotes an identity matrix of dimension $n \in \mathbb{N}$. Clearly,

$$
\begin{equation*}
I_{J P} \preceq \mathcal{E}_{\alpha} \preceq(1+\alpha) I_{J P} \text { and } I_{J P} \succeq \mathcal{E}_{\alpha}^{-1} \succeq(1+\alpha)^{-1} I_{J P}, \tag{44}
\end{equation*}
$$

where, for arbitrary square matrices $X$ and $Y$, we write $X \succeq Y$ if $X-Y \geq 0$, and vice versa. Moreover, from (26), (27), (31) and using the identity $(A \otimes B)(X \otimes Y)=$ $A X \otimes B Y$, we obtain

$$
\begin{align*}
\mathcal{M}_{\alpha} & =G_{0} \otimes M+\alpha T \otimes M \\
& =\left(G_{0}+\alpha T\right) \otimes M \\
& =\left(G_{0} I_{P}+\alpha G_{0} \Upsilon\right) \otimes\left(M I_{J}\right) \\
& =\left(G_{0} \otimes M\right)\left(I_{P} \otimes I_{J}\right)+\left(G_{0} \otimes M\right)\left(\alpha \Upsilon \otimes I_{J}\right) \\
& =\left(G_{0} \otimes M\right)\left[\left(I_{P} \otimes I_{J}\right)+\left(\alpha \Upsilon \otimes I_{J}\right)\right] \\
& =\mathcal{M}\left[\left(I_{P}+\alpha \Upsilon\right) \otimes I_{J}\right] \\
& =\mathcal{M} \mathcal{E}_{\alpha}=\mathcal{E}_{\alpha} \mathcal{M} \tag{45}
\end{align*}
$$

since both $G_{0}$ and $I_{P}+\alpha \Upsilon$ are diagonal matrices and therefore commute with each other. Now, recall from (41) that the approximation $S_{1}$ to the Schur complement $S$ is given by

$$
S_{1}=\left(\mathcal{K} \mathcal{M}_{\alpha}^{-1} \mathcal{K}+\frac{1+\alpha}{\beta} \mathcal{M} \mathcal{M}_{\alpha}^{-1} \mathcal{M}+\sqrt{\frac{1+\alpha}{\beta}}\left[\mathcal{K} \mathcal{M}_{\alpha}^{-1} \mathcal{M}+\mathcal{M} \mathcal{M}_{\alpha}^{-1} \mathcal{K}\right]\right)
$$

so that using (36), (45) and (46), we see that the preconditioned Schur complement $S_{1}^{-1} S$ is similar to the matrix

$$
\begin{equation*}
\mathcal{M}^{1 / 2} S_{1}^{-1} S \mathcal{M}^{-1 / 2}=\left(\mathcal{M}^{-1 / 2} S_{1} \mathcal{M}^{-1 / 2}\right)^{-1}\left(\mathcal{M}^{-1 / 2} S \mathcal{M}^{-1 / 2}\right) \tag{46}
\end{equation*}
$$

It therefore follows that

$$
\begin{aligned}
S_{1}^{-1} S & \sim\left(\mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C}+\frac{1+\alpha}{\beta} \mathcal{E}_{\alpha}^{-1}+\sqrt{\frac{1+\alpha}{\beta}}\left(\mathcal{C E}_{\alpha}^{-1}+\mathcal{E}_{\alpha}^{-1} \mathcal{C}\right)\right)^{-1}\left(\mathcal{C E}_{\alpha}^{-1} \mathcal{C}+\beta^{-1} I_{J P}\right) \\
& =\left(\beta \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C}+(1+\alpha) \mathcal{E}_{\alpha}^{-1}+\sqrt{\beta(1+\alpha)}\left(\mathcal{C E}_{\alpha}^{-1}+\mathcal{E}_{\alpha}^{-1} \mathcal{C}\right)\right)^{-1}\left(\beta \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C}+I_{J P}\right)
\end{aligned}
$$

where $\sim$ implies similarity transformation and $\mathcal{C}:=\mathcal{M}^{-1 / 2} \mathcal{K} \mathcal{M}^{-1 / 2}$. Now, observe that the matrix $\mathcal{C}$ is symmetric and positive definite so that $\lambda(\mathcal{C}) \subset(0,+\infty)$. Consider now the Raleigh quotient

$$
R(x):=\frac{x^{T}\left[\beta \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C}+I_{J P}\right] x}{x^{T}\left[\beta \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C}+(1+\alpha) \mathcal{E}_{\alpha}^{-1}+\sqrt{\beta(1+\alpha)}\left(\mathcal{C E}_{\alpha}^{-1}+\mathcal{E}_{\alpha}^{-1} \mathcal{C}\right)\right] x}
$$

But then, $\kappa\left(\mathcal{E}^{-1}\right)=1+\alpha$, and hence, by Proposition 1. we have that

$$
x^{T}\left(\mathcal{C} \mathcal{E}_{\alpha}^{-1}+\mathcal{E}_{\alpha}^{-1} \mathcal{C}\right) x>0, \text { for } \alpha+1<\left(\frac{\sqrt{\kappa(\mathcal{C})}+1}{\sqrt{\kappa(\mathcal{C})}-1}\right)^{2}
$$

This, in turn, implies that the denominator of $R(x)$ is also strictly positive. Hence,

$$
R(x) \leq \frac{x^{T}\left[\beta \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C}+(1+\alpha) \mathcal{E}_{\alpha}^{-1}\right] x}{x^{T}\left[\beta \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C}+(1+\alpha) \mathcal{E}_{\alpha}^{-1}+\sqrt{\beta(1+\alpha)}\left(\mathcal{C E}_{\alpha}^{-1}+\mathcal{E}_{\alpha}^{-1} \mathcal{C}\right)\right] x}<1
$$

from which we deduce that $\lambda_{\max }:=\max R(x)<1$.
Now, observe that $x^{T} \mathcal{C} \mathcal{E}_{\alpha}^{-1} x=x^{T} \mathcal{E}_{\alpha}^{-1} \mathcal{C} x$. Moreover, for any two vectors $z_{1}, z_{2}$ of appropriate dimensions, Cauchy-Schwarz Inequality implies $\left\langle z_{1}^{T} z_{2}\right\rangle^{2} \leq\left(z_{1}^{T} z_{1}\right)\left(z_{2}^{T} z_{2}\right)$. Thus, setting $z_{1}^{T}=x^{T} \mathcal{C} \mathcal{E}_{\alpha}^{-1 / 2}$ and $z_{2}=\mathcal{E}_{\alpha}^{-1 / 2} x$, we obtain

$$
\begin{equation*}
\left(x^{T} \mathcal{C} \mathcal{E}_{\alpha}^{-1} x\right)^{2} \leq\left(x^{T} \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C} x\right)\left(x^{T} \mathcal{E}_{\alpha}^{-1} x\right) \tag{47}
\end{equation*}
$$

Hence, using 47, together with the fact that $(a+b)^{2} \leq 2\left(a^{2}+b^{2}\right), a, b \in \mathbb{R}$, one gets

$$
\begin{align*}
R(x) & =\frac{x^{T}\left[\beta \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C}+I_{J P}\right] x}{x^{T}\left[\beta \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C}+(1+\alpha) \mathcal{E}_{\alpha}^{-1}+\sqrt{\beta(1+\alpha)}\left(\mathcal{C E}_{\alpha}^{-1}+\mathcal{E}_{\alpha}^{-1} \mathcal{C}\right)\right] x} \\
& \geq \frac{x^{T} \beta \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C} x+x^{T} I_{J P} x}{\left[\beta^{1 / 2}\left(x^{T} \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C} x\right)^{1 / 2}+(1+\alpha)^{1 / 2}\left(x^{T} \mathcal{E}_{\alpha}^{-1} x\right)^{1 / 2}\right]^{2}} \\
& \geq \frac{x^{T} \beta \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C} x+x^{T} I_{J P} x}{2\left[\beta x^{T} \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C} x+(1+\alpha) x^{T} \mathcal{E}_{\alpha}^{-1} x\right]} \\
& \geq \frac{x^{T} \beta \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C} x+x^{T} \mathcal{E}_{\alpha}^{-1} x}{2\left[\beta x^{T} \mathcal{C} \mathcal{E}_{\alpha}^{-1} \mathcal{C} x+(1+\alpha) x^{T} \mathcal{E}_{\alpha}^{-1} x\right]} \\
& \geq \frac{x^{T} \mathcal{E}_{\alpha}^{-1} x}{2(1+\alpha) x^{T} \mathcal{E}_{\alpha}^{-1} x}=\frac{1}{2(1+\alpha)}, \tag{48}
\end{align*}
$$

which shows that $\lambda_{\min }:=\min R(x) \geq \frac{1}{2(1+\alpha)}$, thereby concluding the proof of the theorem.

Note that, in the context of a deterministic optimal control problem, Pearson and Wathen in [24, Theorem 4] have independently obtained, specifically for $\alpha=0$, a
similar result as that of Theorem 3. We, however, point out herein that, in addition to the generalization of the said result, ours yields a sharper bound than the one that these authors obtained. Moreover, with the exception of the parameter $\alpha$, the result of Theorem 3 is independent of the discretization parameters in the system.

The following result is an immediate consequence of Theorem 3
Theorem 4. Let $\mathcal{A}$ be the KKT matrix given by (33) and define $\mathcal{P}_{0}$ by

$$
\mathcal{P}_{0}=\left[\begin{array}{cc}
A & 0 \\
0 & S_{1}
\end{array}\right],
$$

where $A$ and $S_{1}$ are given, respectively, by (34) and 41. Moreover, assume that $\alpha<\left(\frac{\sqrt{\kappa(\mathcal{C})}+1}{\sqrt{\kappa(\mathcal{C})}-1}\right)^{2}-1$, where $\mathcal{C}$ is as defined in Theorem 3 . Then, the eigenvalues of the matrix $\mathcal{P}_{0}^{-1} \mathcal{A}$ satisfy

$$
\begin{equation*}
\lambda\left(\mathcal{P}_{0}^{-1} \mathcal{A}\right)=\{1\} \cup \mathcal{I}^{-} \cup \mathcal{I}^{+} \tag{49}
\end{equation*}
$$

where
$\mathcal{I}^{-}=\left(\frac{1}{2}(1-\sqrt{5}), \frac{1}{2}\left(1-\sqrt{1+\frac{2}{1+\alpha}}\right)\right], \mathcal{I}^{+}=\left[\frac{1}{2}\left(1+\sqrt{1+\frac{2}{1+\alpha}}\right), \frac{1}{2}(1+\sqrt{5})\right)$.
Proof. First, we note that $\mathcal{P}_{0}^{-1} \mathcal{A}$ shares the same eigenvalues with the symmetric matrix given by

$$
\mathcal{P}_{0}^{-1 / 2} \mathcal{A P}_{0}^{-1 / 2}=\left[\begin{array}{cc}
I & A^{-1 / 2} B^{T} S_{1}^{-1 / 2} \\
S_{1}^{-1 / 2} B A^{-1 / 2} & 0
\end{array}\right]
$$

Now, using [10, Lemma 2.1], we know that the eigenvalues of $\mathcal{P}_{0}^{-1 / 2} \mathcal{A} \mathcal{P}_{0}^{-1 / 2}$ are either 1 or have the form $\frac{1}{2}\left(1 \pm \sqrt{1+4 s^{2}}\right)$, where $s$ is a singular value of $X:=S_{1}^{-1 / 2} B A^{-1 / 2}$; in other words, $s^{2}$ is an eigenvalue of $X^{T} X$. Since $S_{1}^{-1} S$ is similar to $X^{T} X$, the result (49) follows immediately from Theorem 3 .

The robustness of $S_{1}$ notwithstanding, we cannot implement it as it is, as this is equivalent to solving the forward problem twice per iteration due to the presence of $\mathcal{Z}:=\mathcal{K}+\mathcal{M}_{u}$ and its transpose. Hence, we need to derive an appropriate approximation for $\mathcal{Z}$. More precisely, observe first, from (26), that

$$
\begin{align*}
\mathcal{Z} & =\mathcal{K}+\sqrt{\frac{1+\alpha}{\beta}} \mathcal{M} \\
& =\left(\sum_{i=0}^{N} G_{i} \otimes K_{i}\right)+\sqrt{\frac{1+\alpha}{\beta}}\left(G_{0} \otimes M\right)=\sum_{i=0}^{N} G_{i} \otimes \tilde{K}_{i}, \tag{50}
\end{align*}
$$

with $\tilde{K}_{0}:=K_{0}+\sqrt{\frac{1+\alpha}{\beta}} M, \quad \tilde{K}_{i}=K_{i}, i=1, \ldots, N$. But then, we can now approximate $\mathcal{Z}$ using similar preconditioners for the stationary forward problems considered in, for example, [25, 32]:

$$
\begin{equation*}
\mathcal{Z}_{0}:=G_{0} \otimes \tilde{K}_{0} \tag{51}
\end{equation*}
$$

or

$$
\begin{equation*}
\mathcal{Z}_{1}:=\sum_{i=0}^{N} \frac{\operatorname{trace}\left(\tilde{K}_{i}^{T} \tilde{K}_{0}\right)}{\operatorname{trace}\left(\tilde{K}_{0}^{T} \tilde{K}_{0}\right)} G_{i} \otimes \tilde{K}_{0} \tag{52}
\end{equation*}
$$

Note that for a practical algorithm, $S_{1}$ is approximated using multigrid techniques for $\tilde{K}_{0}$ in both $\mathcal{Z}_{0}$ and $\mathcal{Z}_{1}$. The preconditioner (51) is the so-called mean-based preconditioner. It is block diagonal and is best suited for systems for which the variance of the random input is small relative to its mean. Its performance deteriorates with increasing variance. This is quite intuitive since from $(24),(25)$ and (50), we can see that as $\sigma_{\kappa}$ increases, the off-diagonal blocks of the global stochastic Galerkin matrix $\sum_{i=1}^{N} G_{i} \otimes K_{i}$ become more significant and they are not represented in the preconditioner. The latter was proposed by Ullmann in [32] to circumvent the shortcomings of the former in the forward problem. Both of them have been successfully applied to a time-dependent forward problem considered in [2]. It is, however, more expensive to implement $\mathcal{Z}_{1}$ than the mean-based preconditioner. Hence, in our numerical experiments, we shall stick to the mean-based preconditioner.

In a nutshell, we outline below the dominant operations in the application of our proposed block-diagonal preconditioner $\mathcal{P}$ in (35).

- $(\mathbf{1}, \mathbf{1}): 1$ Chebyshev semi-iteration for the mass matrix $M$.
- $(\mathbf{2 , 2}): 1$ Chebyshev semi-iteration for the mass matrix $M$.
- $(\mathbf{3}, \mathbf{3}): 2$ multigrid operations: 1 for $\mathcal{Z}_{0}$ and 1 for its transpose.
- Total: 2 Chebyshev semi-iterations and 2 multigrid operations.

Having been equipped with a suitable preconditioner, we proceed to the next section to discuss our Krylov subspace solver.

### 2.2 Computing low-rank approximation of the solution to the stationary problem

As we have already pointed out in Section 2.1, the MINRES algorithm is an optimal solver for the system (30). Hence, we will use it, together with (35), to solve 30). In particular, our approach is based on the low-rank version of MINRES presented in [28. In this section, we give a brief overview of this low-rank iterative solver. Now, observe first that using the identity

$$
\begin{equation*}
\operatorname{vec}(W X V)=\left(V^{T} \otimes W\right) \operatorname{vec}(X) \tag{53}
\end{equation*}
$$

where

$$
\operatorname{vec}(X)=\left(\begin{array}{c}
x_{1} \\
\vdots \\
x_{m}
\end{array}\right) \in \mathbb{R}^{m \times 1}
$$

for any $X=\left[x_{1}, \ldots, x_{m}\right] \in \mathbb{R}^{n \times m}$, the linear system can be rewritten as $\mathcal{A} \mathcal{X}=\mathcal{R}$, where

$$
\begin{gathered}
\mathcal{A}=\left[\begin{array}{ccc}
G_{\alpha} \otimes M & 0 & -\sum_{i=0}^{N} G_{i} \otimes K_{i} \\
0 & \beta\left(G_{0} \otimes M\right) & G_{0} \otimes M \\
-\sum_{i=0}^{N} G_{i} \otimes K_{i} & G_{0} \otimes M & 0
\end{array}\right] \\
\mathcal{X}=\left[\begin{array}{c}
\operatorname{vec}(Y) \\
\operatorname{vec}(U) \\
\operatorname{vec}(F)
\end{array}\right], \quad \mathcal{R}=\left[\begin{array}{c}
\operatorname{vec}\left(R_{1}\right) \\
0 \\
\operatorname{vec}\left(R_{3}\right)
\end{array}\right]
\end{gathered}
$$

and

$$
\begin{gathered}
Y=\left[y_{0}, \ldots, y_{P-1}\right], \quad U=\left[u_{0}, \ldots, u_{P-1}\right], \quad F=\left[f_{0}, \ldots, f_{P-1}\right] \\
R_{1}=\operatorname{vec}^{-1}\left(\left(G_{0} \otimes M\right) \overline{\mathbf{y}}\right), \quad R_{3}=\operatorname{vec}^{-1}(\mathbf{d})
\end{gathered}
$$

Hence, (53) implies that

$$
\mathcal{A X}=\operatorname{vec}\left(\left[\begin{array}{c}
M Y G_{\alpha}^{T}-\sum_{i=0}^{N} K_{i} F G_{i}^{T}  \tag{54}\\
\beta M U G_{0}^{T}+M F G_{0}^{T} \\
-\sum_{i=0}^{N} K_{i} Y G_{i}^{T}+M U G_{0}^{T}
\end{array}\right]\right)=\operatorname{vec}\left(\left[\begin{array}{c}
R_{1} \\
0 \\
R_{3}
\end{array}\right]\right) .
$$

Our approach is essentially based on the assumption that both the solution matrix $\mathcal{X}$ and the right hand side matrix $\mathcal{R}$ admit low-rank representations; that is,

$$
\left\{\begin{array}{lll}
Y=W_{Y} V_{Y}^{T}, & \text { with } W_{Y} \in \mathbb{R}^{J \times k_{1}}, & V_{Y} \in \mathbb{R}^{P \times k_{1}}  \tag{55}\\
U=W_{U} V_{U}^{T}, & \text { with } W_{U} \in \mathbb{R}^{J \times k_{2}}, & V_{U} \in \mathbb{R}^{P \times k_{2}} \\
F=W_{F} V_{F}^{T}, & \text { with } W_{F} \in \mathbb{R}^{J \times k_{3}}, & V_{F} \in \mathbb{R}^{P \times k_{3}},
\end{array}\right.
$$

where $k_{1,2,3}$ are small relative to $P$. Substituting (55) in and ignoring the vec operator, we then obtain

$$
\left[\begin{array}{c}
M W_{Y} V_{Y}^{T} G_{\alpha}^{T}-\sum_{i=0}^{N} K_{i} W_{F} V_{F}^{T} G_{i}^{T}  \tag{56}\\
\beta M W_{U} V_{U}^{T} G_{0}^{T}+M W_{F} V_{F}^{T} G_{0}^{T} \\
-\sum_{i=0}^{N} K_{i} W_{Y} V_{Y}^{T} G_{i}^{T}+M W_{U} V_{U}^{T} G_{0}^{T}
\end{array}\right]=\left[\begin{array}{c}
R_{11} R_{12}^{T} \\
0 \\
R_{31} R_{32}^{T}
\end{array}\right]
$$

where $R_{11} R_{12}^{T}$ and $R_{31} R_{32}^{T}$ are the low-rank representations of the $R_{1}$ and $R_{3}$, respectively.

The attractiveness of this approach lies therefore in the fact that one can rewrite the three block rows in the left hand side in (56), respectively, as
so that the low-rank nature of the factors guarantees fewer multiplications with the submatrices while maintaining smaller storage requirements. More precisely, keeping in mind that

$$
x=\operatorname{vec}\left(\left[\begin{array}{c}
X_{11} X_{12}^{T} \\
X_{21} X_{22}^{T} \\
X_{31} X_{32}^{T}
\end{array}\right]\right)
$$

corresponds to the associated vector $x$ from a vector-based version of MINRES, matrixvector multiplication in our low-rank MINRES is given by Algorithm 2. Note that an

```
Algorithm 2 Matrix-vector multiplication in low-rank MINRES
    Input: \(W_{11}, W_{12}, W_{21}, W_{22}, W_{31}, X_{32}\)
    Output: \(X_{11}, X_{12}, X_{21}, X_{22}, X_{31}, X_{32}\)
    \(X_{11}=\left[\begin{array}{ll}M W_{11} & -\sum_{i=0}^{N} K_{i} W_{31}\end{array}\right]\)
    \(X_{12}=\left[\begin{array}{ll}G_{\alpha} W_{12} & G_{i} W_{32}\end{array}\right]\)
    \(X_{21}=\left[\begin{array}{ll}\beta G_{0} W_{21} & -\sum_{i=0}^{N} K_{i} W_{31}\end{array}\right]\)
    \(X_{22}=\left[\begin{array}{ll}G_{0} W_{22} & G_{0} W_{32}\end{array}\right]\)
    \(X_{31}=\left[-\sum_{i=0}^{N} K_{i} W_{11} \quad M W_{21}\right]\)
    \(X_{32}=\left[\begin{array}{ll}G_{i} W_{12} & G_{0} W_{22}\end{array}\right]\)
```

important feature of low-rank MINRES is that the iterates of the solution matrices $Y, U$ and $F$ in the algorithm are truncated by a truncation operator $\mathcal{T}_{\epsilon}$ with a prescribed tolerance $\epsilon$. This is accomplished via QR decomposition as in [16] or truncated singular value decomposition (SVD) as in [2, 28.
The truncation operation is necessary because the new computed factors could have increased ranks compared to the original factors in (57). Hence, a truncation of all the
factors after the matrix-vector products, is used to construct new factors; for instance,

$$
\left[\tilde{X}_{11}, \tilde{X}_{12}\right]:=\mathcal{T}_{\epsilon}\left(\left[X_{11}, X_{12}\right]\right)=\mathcal{T}_{\epsilon}\left(\left[\begin{array}{ll}
M W_{11} & -\sum_{i=0}^{N} K_{i} W_{31}
\end{array}\right]\left[\begin{array}{c}
W_{12}^{T} G_{\alpha}^{T} \\
W_{32}^{T} G_{i}^{T}
\end{array}\right]\right)
$$

Moreover, in order to ensure that the inner products within the iterative low-rank solver are computed efficiently, we use the fact that

$$
\langle x, y\rangle=\operatorname{vec}(X)^{T} \operatorname{vec}(Y)=\operatorname{trace}\left(X^{T} Y\right)
$$

to deduce that

$$
\operatorname{trace}(\underbrace{\left(X_{1} X_{2}^{T}\right)^{T}}_{\text {Large }} \underbrace{\left(Y_{1} Y_{2}^{T}\right)}_{\text {Large }})=\operatorname{trace}(\underbrace{Y_{2}^{T} X_{2}}_{\text {Small }} \underbrace{X_{1}^{T} Y_{1}}_{\text {Small }}),
$$

where $X=X_{1} X_{2}^{T}$ and $Y=Y_{1} Y_{2}^{T}$, which allows us to compute the trace of small matrices rather than of the ones from the full model.
For more details on implementation issues, we refer the interested reader to [2, 28]. In Section 4, we use numerical experiments to illustrate the performance of low-rank MINRES, together with the preconditioners discussed in Section 2.1.

Next, we proceed to Section 3 to present a time-dependent analogue of the model problem considered so far.

## 3 A stochastic parabolic optimal control problem

In an attempt to extend our discussion on the above model problem to a timedependent case, we henceforth replace $L^{2}(\mathcal{D})$ in (7) by the space

$$
L^{2}([0, T], \mathcal{D})=\left\{f \in L^{2}(\mathcal{D}): \int_{0}^{T}[f(t)]^{2} d t<\infty\right\}
$$

and then consider a parabolic SOCP now given by $\mathcal{J}(y(t, \mathbf{x}, \omega)), u(t, \mathbf{x}, \omega))$ subject, $\mathbb{P}$-almost surely, to

$$
\left\{\begin{align*}
\frac{\partial y(t, \mathbf{x}, \omega)}{\partial t}-\nabla \cdot(a(\mathbf{x}, \omega) \nabla y(t, \mathbf{x}, \omega)) & =u(t, \mathbf{x}, \omega), \text { in }(0, T] \times \mathcal{D} \times \Omega  \tag{58}\\
y(t, \mathbf{x}, \omega) & =0, \text { on }(0, T] \times \partial \mathcal{D} \times \Omega \\
y(0, \mathbf{x}, \omega) & =y_{0}, \quad \text { in } \mathcal{D} \times \Omega
\end{align*}\right.
$$

where the random control function satisfies

$$
\int_{\Omega} u(\cdot, \cdot, \omega) d \mathbb{P}(\omega)<+\infty, \text { a.e }
$$

and, as before, $a(\mathbf{x}, \omega)$ is assumed to be uniformly positive in $\mathcal{D} \times \Omega$. We note here that the time-dependence of this problem introduces an additional degree of freedom which
makes the system matrix here (a lot) larger than the system matrix in the steady-state case.
We use the trapezoidal rule for temporal discretization (as was done for deterministic problems in e.g. [23, 28]) and SGFEM in the spatial and the stochastic domains to get the following discrete objective function

$$
\begin{equation*}
\mathcal{J}(y, u)=\frac{\tau}{2}(\mathbf{y}-\overline{\mathbf{y}})^{T} \mathcal{M}_{a}(\mathbf{y}-\overline{\mathbf{y}})+\frac{\tau \alpha}{2} \mathbf{y}^{T} \mathcal{M}_{b} \mathbf{y}+\frac{\tau \beta}{2} \mathbf{u}^{T} \mathcal{M}_{2} \mathbf{u} \tag{59}
\end{equation*}
$$

where

$$
\left\{\begin{array}{l}
\mathcal{M}_{a}=\operatorname{blkdiag}\left(\frac{1}{2} \mathcal{M}, \mathcal{M}, \ldots, \mathcal{M}, \frac{1}{2} \mathcal{M}\right)  \tag{60}\\
\mathcal{M}_{b}=\operatorname{blkdiag}\left(\frac{1}{2} \mathcal{M}_{t}, \mathcal{M}_{t}, \ldots, \mathcal{M}_{t}, \frac{1}{2} \mathcal{M}_{t}\right)
\end{array}\right.
$$

with $\mathcal{M}$ and $\mathcal{M}_{t}$ as defined in 26) and respectively. Note that $\mathcal{M}_{2}=\mathcal{M}_{a}$. Here, denoting the number of time steps by $n_{t}$, we also note that

$$
\mathbf{y}=\left[\begin{array}{c}
\mathbf{y}_{1} \\
\vdots \\
\mathbf{y}_{n_{t}}
\end{array}\right], \overline{\mathbf{y}}=\left[\begin{array}{c}
\overline{\mathbf{y}}_{1} \\
\vdots \\
\overline{\mathbf{y}}_{n_{t}}
\end{array}\right] \text { and } \mathbf{u}=\left[\begin{array}{c}
\mathbf{u}_{1} \\
\vdots \\
\mathbf{u}_{n_{t}}
\end{array}\right]
$$

with $\mathbf{y}_{i}, \overline{\mathbf{y}}_{i}, \mathbf{u}_{i} \in \mathbb{R}^{J P \times 1}, \quad i=1, \ldots, n_{t}$.
For an all-at-once discretization of the state equation (58), we use the implicit Euler method together with SGFEM to get

$$
\mathcal{K}_{t} \mathbf{y}-\tau \mathcal{N} \mathbf{u}=\mathbf{d}
$$

where

$$
\mathcal{K}_{t}=\left[\begin{array}{cccc}
\mathcal{L} & & & \\
-\mathcal{M} & \mathcal{L} & & \\
& \ddots & \ddots & \\
& & -\mathcal{M} & \mathcal{L}
\end{array}\right], \mathcal{N}=\left[\begin{array}{cccc}
\mathcal{M} & & & \\
& \mathcal{M} & & \\
& & \ddots & \\
& & & \mathcal{M}
\end{array}\right], \mathbf{d}=\left[\begin{array}{c}
\mathcal{M} \mathbf{y}_{0} \\
0 \\
\vdots \\
0
\end{array}\right]
$$

where $\mathcal{L}:=G_{0} \otimes\left(M+\tau K_{0}\right)+\tau \sum_{i=1}^{N} G_{i} \otimes K_{i}$. Observe that the matrix $\mathcal{K}_{t}$ in this case is not symmetric, unlike the matrix $\mathcal{K}$ in the stationary case.

Applying first order conditions to the Lagrangian functional for this constrained optimization problem yields

$$
\left[\begin{array}{ccc}
\tau \mathcal{M}_{1} & 0 & -\mathcal{K}_{t}^{T}  \tag{61}\\
0 & \beta \tau \mathcal{M}_{2} & \tau \mathcal{N}^{T} \\
-\mathcal{K}_{t} & \tau \mathcal{N} & 0
\end{array}\right]\left[\begin{array}{l}
\mathbf{y} \\
\mathbf{u} \\
\mathbf{f}
\end{array}\right]=\left[\begin{array}{c}
\tau \mathcal{M}_{a} \overline{\mathbf{y}} \\
\mathbf{0} \\
\mathbf{d}
\end{array}\right]
$$

where, from (60) and 66),

$$
\begin{align*}
\mathcal{M}_{1} & =\mathcal{M}_{a}+\alpha \mathcal{M}_{b} \\
& =(D \otimes \mathcal{M})+\alpha\left(D \otimes \mathcal{M}_{t}\right)=D \otimes G_{\alpha} \otimes M=D \otimes \mathcal{M}_{\alpha} \tag{62}
\end{align*}
$$

with $G_{\alpha}$ as defined in (32), and

$$
\begin{equation*}
D=\operatorname{diag}\left(\frac{1}{2}, 1 \ldots, 1, \frac{1}{2}\right) \in \mathbb{R}^{n_{t} \times n_{t}} \tag{63}
\end{equation*}
$$

We note here that

$$
\begin{equation*}
\mathcal{K}_{t}=\left(I_{n_{t}} \otimes \mathcal{L}\right)+(C \otimes \mathcal{M})=I_{n_{t}} \otimes\left[\sum_{i=0}^{N} G_{i} \otimes \hat{K}_{i}\right]+\left(C \otimes G_{0} \otimes M\right), \tag{64}
\end{equation*}
$$

where $\hat{K}_{0}=M+\tau K_{0}, \hat{K}_{i}=\tau K_{i}, i=1, \ldots, N$. The matrix $C \in \mathbb{R}^{n_{t} \times n_{t}}$ comes from the implicit Euler discretization and is given by

$$
C=\left[\begin{array}{cccc}
0 & & & \\
-1 & 0 & & \\
& \ddots & \ddots & \\
& & -1 & 0
\end{array}\right]
$$

and $I_{n_{t}}$ is an identity matrix of dimension $n_{t}$. The use of other temporal discretizations is, of course, possible. The Crank-Nicholson scheme, for instance, can be written in a similar way. Moreover,

$$
\begin{equation*}
\mathcal{N}=I_{n_{t}} \otimes G_{0} \otimes M, \quad \mathcal{M}_{2}=D \otimes G_{0} \otimes M \tag{65}
\end{equation*}
$$

Hence, each of the block matrices $\mathcal{K}_{t}, \mathcal{N}, \mathcal{M}_{1}$ and $\mathcal{M}_{2}$ belongs to $\mathbb{R}^{J P n_{t} \times J P n_{t}}$, since $G_{i} \in \mathbb{R}^{P \times P}, i=0, \ldots, P-1$, and $M, K_{i} \in \mathbb{R}^{J \times J}, i=0, \ldots, N$. So, the overall coefficient matrix in (61) is of dimension $3 J P n_{t} \times 3 J P n_{t}$.

As can be seen from (64), for instance, the time-dependent problem leads to an additional Kronecker product. Indeed, although the low-rank solver presented in the stationary case reduces storage problems in large-scale simulations, the low-rank factors become infeasible in higher dimensions. Further data compression can, fortunately, be achieved with more advanced high-dimensional tensor product decompositions. Together with preconditioned MINRES, we henceforth solve the linear system discussed in this section using an elegant and robust tensor format called Tensor Train (TT) format which was introduced in [20]. To that end, we proceed next to Section 3.1 to give an overview of the TT format.

### 3.1 Solving the optimality systems from the unsteady problem

First, we recall that a tensor $\mathbf{y}:=\mathbf{y}\left(i_{1}, \ldots, i_{d}\right), i_{k}=1, \ldots, n_{k}$ is an $n_{1} \times n_{2} \times \ldots \times n_{d}$ multi-dimensional array, where the integers $n_{1}, n_{2}, \ldots, n_{d}$ are called the mode sizes and $d$ is the order of $\mathbf{y}$. The tensor $\mathbf{y}$ admits a tensor train decomposition or TTformat [20, 7] if it can be expressed as

$$
\mathbf{y}\left(i_{1}, \ldots, i_{d}\right)=\mathbf{y}_{1}\left(i_{1}\right) \mathbf{y}_{2}\left(i_{2}\right) \cdots \mathbf{y}_{d}\left(i_{d}\right)
$$

where $\mathbf{y}_{k}\left(i_{k}\right)$ is an $r_{k-1} \times r_{k}$ matrix for each fixed $i_{k}, 1 \leq i_{k} \leq n_{k}$. Moreover, the numbers $r_{k}$ are called the TT-ranks, whereas $\mathbf{y}_{k}\left(i_{k}\right)$ are the cores of the decomposition. More precisely, $\mathbf{y}_{k}\left(i_{k}\right)$ is a three-dimensional array, and it can essentially be treated as an $r_{k-1} \times n_{k} \times r_{k}$ array with elements $\mathbf{y}_{k}\left(\alpha_{k-1}, n_{k}, \alpha_{k}\right)=\mathbf{y}_{\alpha_{k-1}, \alpha_{k}}^{(k)}\left(i_{k}\right)$. Here, the boundary conditions $r_{0}=r_{d}=1$ are imposed on the decomposition to make the matrix-by-matrix products a scalar. The decomposition can be expressed in index form as

$$
\begin{equation*}
\mathbf{y}\left(i_{1}, \ldots, i_{d}\right)=\sum_{\alpha_{1} \ldots \alpha_{d-1}=1}^{r_{1} \ldots r_{d-1}} \mathbf{y}_{1}\left(\alpha_{0}, n_{1}, \alpha_{1}\right) \mathbf{y}_{2}\left(\alpha_{1}, n_{2}, \alpha_{2}\right) \cdots \mathbf{y}_{d-1}\left(\alpha_{d-1}, n_{d}, \alpha_{d}\right) \tag{66}
\end{equation*}
$$

where $\alpha_{0}=\alpha_{d}=1$. It turns out that TT-decomposition yields a low-rank format for tensors as it is derived by a repeated application of low-rank approximation 20. To see this [8, let

$$
\begin{equation*}
\overline{i_{2} \cdots i_{d}}=i_{2}+\left(i_{3}-1\right) n_{2}+\cdots+\left(i_{d}-1\right) n_{2} n_{3} \cdots n_{d-1} . \tag{67}
\end{equation*}
$$

Then, by regrouping of indices, one can rewrite $\mathbf{y}$ as a matrix $Y_{1} \in \mathbb{R}^{n_{1} \times n_{2} \cdots n_{d}}$ with $Y_{1}\left(i_{1}, \overline{i_{2} \cdots i_{d}}\right)=\mathbf{y}\left(i_{1}, \ldots, i_{d}\right)$. Thus, applying a low-rank SVD to the matrix $Y_{1}$ yields

$$
Y_{1} \approx U_{1} \Sigma_{1} V_{1}^{T}, \quad U_{1} \in \mathbb{R}^{n_{1} \times r_{1}}, \quad V_{1} \in \mathbb{R}^{n_{2} \cdots n_{d} \times r_{1}}
$$

The first factor $U_{1}$ is of moderate dimension and can be stored as $\mathbf{y}_{\alpha_{1}}^{(1)}\left(i_{1}\right)=U_{1}\left(i_{1}, \alpha_{1}\right)$, where $\alpha_{1}=1, \ldots, r_{1}$. The remaining matrix $\Sigma_{1} V_{1}^{T}$ depends on $\alpha_{1}$ and $i_{2} \cdots i_{d}$. Next, we regroup these indices as follows

$$
Y_{2}\left(\overline{\alpha_{1} i_{2}}, \overline{i_{3} \cdots i_{d}}\right)=\Sigma_{1}\left(\alpha_{1}, \alpha_{1}\right) V_{1}^{T}\left(\alpha_{1}, \overline{i_{2} \cdots i_{d}}\right)
$$

and compute the next SVD:

$$
Y_{2} \approx U_{2} \Sigma_{2} V_{2}^{T}, \quad U_{2} \in \mathbb{R}^{r_{1} n_{2} \times r_{2}}, \quad V_{2} \in \mathbb{R}^{n_{3} \cdots n_{d} \times r_{2}} .
$$

Again, $U_{2}$ can be reshaped to a 3 D tensor $\mathbf{y}_{\alpha_{1}, \alpha_{2}}^{(2)}\left(i_{2}\right)=U_{2}\left(\overline{\alpha_{1} i_{2}}, \alpha_{2}\right)$ of moderate size, and the decomposition also applied to $\Sigma_{2} V_{2}^{T}$. Proceeding in this manner, one eventually obtains the TT format:

$$
\begin{equation*}
\mathbf{y}\left(i_{1}, \ldots, i_{d}\right)=\sum_{\alpha_{1} \ldots \alpha_{d-1}=1}^{r_{1} \ldots r_{d-1}} \mathbf{y}_{\alpha_{1}}^{(1)}\left(i_{1}\right) \mathbf{y}_{\alpha_{1}, \alpha_{2}}^{(2)}\left(i_{2}\right) \cdots \mathbf{y}_{\alpha_{d-2}, \alpha_{d-1}}^{(d-1)}\left(i_{d-1}\right) \mathbf{y}_{\alpha_{d-1}}^{(d)}\left(i_{d}\right), \tag{68}
\end{equation*}
$$

with the total storage of at most $d n r^{2}$ memory cells, where $r_{k} \leq r, n_{k} \leq n$. In particular, if $r$ is small, then this requirement is much smaller than the storage of the full array, $n^{d}$. A similar construction can be made for discretized operators in high dimensions. To this end, consider a matrix $A=A\left(\overline{i_{1} \cdots i_{d}}, \overline{j_{1} \cdots j_{d}}\right) \in \mathbb{R}^{\left(n_{1} \cdots n_{d}\right) \times\left(n_{1} \cdots n_{d}\right)}$. We decompose $A$ as follows:

$$
\begin{equation*}
A\left(\overline{i_{1} \cdots i_{d}}, \overline{j_{1} \cdots j_{d}}\right) \approx \sum_{\beta_{1} \ldots \beta_{d-1}=1}^{R_{1} \ldots R_{d-1}} \mathbf{A}_{\beta_{1}}^{(1)}\left(i_{1}, j_{1}\right) \mathbf{A}_{\beta_{1}, \beta_{2}}^{(2)}\left(i_{2}, j_{2}\right) \cdots \mathbf{A}_{\beta_{d-1}}^{(d)}\left(i_{d}, j_{d}\right), \tag{69}
\end{equation*}
$$

which is consistent with the Kronecker product $A=A^{(1)} \otimes A^{(2)}$ in the case $d=2$ and $R_{1}=1$, and allows a natural multiplication with returning the result in the same format.

As pointed out in [7, the TT-format is stable in the sense that one can always find the best approximation of tensors computed via a sequence of QR and SVD decompositions of auxiliary matrices. The TT-decomposition algorithm is implemented in the TT-toolbox [21] and comes with a number of basic linear algebra operations, such as addition, subtraction, matrix-by-vector product, etc. Unfortunately, these operations lead to prohibitive increase in the TT-ranks. Thus, one necessarily has to truncate (or round) the resulting tensor after implementing each of the operations. This enhances the efficiency of the method when used with any standard iterative method such as MINRES. We point out that although solving the KKT system in TT-format (and, in general, with low-rank solvers) introduces further error in the simulation due to the low-rank truncations, the relative tolerance of the truncation operator can be so tightened that the error will become negligible. This is investigated in [2] for a low-rank conjugate gradient iterative solver; see also e.g. [7, 8] for TT iterative solvers.

We remark here that there are, of course, other tensor formats such as canonical, hierarchical and Tucker formats which could be used to represent tensors 12 and hence solve our linear systems. However, our choice of TT-format (or TT toolbox) is due to its relative elegance and convenience in implementation. The details of its implementation are found in [21]. A comprehensive overview of low-rank tensor decompositions can be found in 12 and the references therein. In our numerical experiments, we use preconditioned MINRES, together with the TT toolbox, to solve the linear system (61).

### 3.2 Preconditioning the optimality system

As in the case of the optimality system associated with the stationary model problem, we need a good preconditioner to solve 61). To this end, we will proceed as before and rewrite the saddle point system (61) as

$$
A=\left[\begin{array}{cc}
\tau \mathcal{M}_{1} & 0  \tag{70}\\
0 & \tau \beta \mathcal{M}_{2}
\end{array}\right], \quad B=\left[\begin{array}{ll}
-\mathcal{K}_{t} & \tau \mathcal{N}
\end{array}\right]
$$

in the notation of (33). Next, we are interested in a block diagonal preconditioner to approximate 61. More precisely, we seek a preconditioner of the form

$$
\hat{\mathcal{P}}=\left[\begin{array}{lll}
A_{1} & & \\
& A_{2} & \\
& & \hat{S}
\end{array}\right],
$$

where the blocks $A_{1} \approx \tau D \otimes G_{\alpha} \otimes M$ and $A_{2} \approx \tau \beta D \otimes G_{0} \otimes M$, and as we noted before, both approximations could be accomplished by applying a Chebyshev semi-iteration on the mass matrix $M$ in the blocks. The matrices $D, G_{0}$ and $G_{\alpha}$ are easy to invert since they are diagonal matrices. Moreover, $\hat{S}$ is an approximation to the (negative)

Schur complement $S_{t}=B A^{-1} B^{T}$, that is,

$$
\begin{equation*}
S_{t}:=\frac{1}{\tau} \mathcal{K}_{t} \mathcal{M}_{1}^{-1} \mathcal{K}_{t}^{T}+\frac{\tau}{\beta} \mathcal{N} \mathcal{M}_{2}^{-1} \mathcal{N}^{T} \tag{71}
\end{equation*}
$$

As in the time-independent case, we consider the following approximation of the Schur complement:

$$
\begin{equation*}
\hat{S}=\frac{1}{\tau}\left(\mathcal{K}_{t}+\hat{\mathcal{M}}_{u}\right) \mathcal{M}_{1}^{-1}\left(\mathcal{K}_{t}+\hat{\mathcal{M}}_{u}\right)^{T} \tag{72}
\end{equation*}
$$

where $\hat{\mathcal{M}}_{u}$ is again determined via the 'terms-matching' procedure so that both the first and second terms in $S_{t}$ and $\hat{S}$ are matched, but the cross terms in $\hat{S}$ are ignored; that is, we have

$$
\hat{\mathcal{M}}_{u} \mathcal{M}_{1}^{-1} \hat{\mathcal{M}}_{u}=\frac{\tau^{2}}{\beta} \mathcal{N} \mathcal{M}_{2}^{-1} \mathcal{N}^{T}
$$

from which we deduce that $\hat{\mathcal{M}}_{u}=\gamma \mathcal{N}$, with $\gamma=\tau \sqrt{\frac{1+\alpha}{\beta}}$, by using a similar arguments as before, so that

$$
\begin{equation*}
\hat{S}=\frac{1}{\tau} \underbrace{\left(\mathcal{K}_{t}+\tau \sqrt{\frac{1+\alpha}{\beta}} \mathcal{N}\right)}_{:=\hat{\mathcal{Z}}} \mathcal{M}_{1}^{-1}\left(\mathcal{K}_{t}+\tau \sqrt{\frac{1+\alpha}{\beta}} \mathcal{N}\right)^{T} \tag{73}
\end{equation*}
$$

Moreover, as in the stationary case, we have the following result regarding the eigenvalues of the preconditioned Schur complement $\hat{S}^{-1} S_{t}$.

Theorem 5. Let $\alpha \in[0,+\infty)$. Then, the eigenvalues of $\hat{S}^{-1} S_{t}$ satisfy

$$
\begin{equation*}
\lambda\left(\hat{S}^{-1} S_{t}\right) \subset\left[\frac{1}{2(1+\alpha)}, 1\right), \quad \forall \alpha<\left(\frac{\sqrt{\kappa(\mathcal{K})}+1}{\sqrt{\kappa(\mathcal{K})}-1}\right)^{2}-1 \tag{74}
\end{equation*}
$$

where $\mathcal{K}=\sum_{i=0}^{N} G_{i} \otimes K_{i}$.
Proof. Let $I_{n_{t}}:=I$, and observe first from (64) that we can rewrite $\mathcal{K}_{t}$ as

$$
\begin{equation*}
\mathcal{K}_{t}=(I+C) \otimes\left(G_{0} \otimes M\right)+I \otimes \tau \sum_{i=0}^{N}\left(G_{i} \otimes K_{i}\right)=J_{0} \otimes \mathcal{M}+\tau I \otimes \mathcal{K} \tag{75}
\end{equation*}
$$

where

$$
J_{0}=I+C=\left[\begin{array}{cccc}
1 & & & \\
-1 & 1 & & \\
& \ddots & \ddots & \\
& & -1 & 1
\end{array}\right]
$$

and $\mathcal{K}$, the coefficient matrix of the stationary forward problem, is positive definite. Now, using (43), (45), (62), (65), we see that

$$
\begin{aligned}
\mathcal{M}_{1} & =D \otimes \mathcal{M}_{\alpha} \\
& =D \otimes \mathcal{M} \mathcal{E}_{\alpha} \\
& =(D \otimes \mathcal{M})\left(I \otimes \mathcal{E}_{\alpha}\right) \\
& =\mathcal{M}_{2} \mathcal{F}_{\alpha}=\mathcal{F}_{\alpha} \mathcal{M}_{2},
\end{aligned}
$$

where $\mathcal{F}_{\alpha}=I \otimes \mathcal{E}_{\alpha}$. Next, define the matrix $\mathcal{X}$ by

$$
\mathcal{X}:=(D \otimes I) \mathcal{M}_{2}^{-1 / 2} \mathcal{K} \mathcal{M}_{2}^{-1 / 2}=D^{1 / 2} J_{0} D^{-1 / 2} \otimes I+\tau I \otimes \mathcal{M}^{-1 / 2} \mathcal{K} \mathcal{M}^{-1 / 2}
$$

Note then that $\mathcal{X}$ is similar to $J_{0} \otimes I+\tau I \otimes \mathcal{M}^{-1} \mathcal{K}$. Moreover, from (71) and (73), we see that $\hat{S}^{-1} S_{t}$ is similar to

$$
\begin{array}{r}
{\left[(D \otimes I) \mathcal{M}_{2}^{-1 / 2} \hat{S} \mathcal{M}_{2}^{-1 / 2}(D \otimes I)\right]^{-1}\left[(D \otimes I) \mathcal{M}_{2}^{-1 / 2} S_{t} \mathcal{M}_{2}^{-1 / 2}(D \otimes I)\right]=} \\
{\left[\beta \mathcal{X} \mathcal{F}_{\alpha}^{-1} \mathcal{X}^{T}+\tau^{2}(1+\alpha) \mathcal{F}_{\alpha}^{-1}+\sqrt{\beta(1+\alpha)}\left(\mathcal{X} \mathcal{F}_{\alpha}^{-1}+\mathcal{F}_{\alpha}^{-1} \mathcal{X}^{T}\right)\right]^{-1}\left(\beta \mathcal{X} \mathcal{F}_{\alpha}^{-1} \mathcal{X}^{T}+\tau^{2} I\right)}
\end{array}
$$

Now, consider the Raleigh quotient

$$
R(x):=\frac{x^{T}\left[\beta \mathcal{X} \mathcal{F}_{\alpha}^{-1} \mathcal{X}^{T}+\tau^{2} I\right] x}{x^{T}\left[\beta \mathcal{X} \mathcal{F}_{\alpha}^{-1} \mathcal{X}^{T}+(1+\alpha) \mathcal{F}_{\alpha}^{-1}+\sqrt{\beta(1+\alpha)}\left(\mathcal{X} \mathcal{F}_{\alpha}^{-1}+\mathcal{F}_{\alpha}^{-1} \mathcal{X}^{T}\right)\right] x}
$$

But then,
$\mathcal{X} \mathcal{F}_{\alpha}^{-1}+\mathcal{F}_{\alpha}^{-1} \mathcal{X}^{T}=D^{1 / 2}\left(J_{0}+J_{0}^{T}\right) D^{-1 / 2} \otimes \mathcal{E}_{\alpha}^{-1}+\tau I \otimes \mathcal{M}^{-1 / 2}\left(\mathcal{K} \mathcal{E}_{\alpha}^{-1}+\mathcal{E}_{\alpha}^{-1} \mathcal{K}\right) \mathcal{M}^{-1 / 2}$.
Since the matrix $D^{1 / 2}\left(J_{0}+J_{0}^{T}\right) D^{-1 / 2}$ is the sum of two positive definite matrices, it is therefore positive definite. Besides, by Proposition 1. one gets

$$
\mathcal{K} \mathcal{E}_{\alpha}^{-1}+\mathcal{E}_{\alpha}^{-1} \mathcal{K} \succ 0, \quad \forall \alpha<\left(\frac{\sqrt{\kappa(\mathcal{K})}+1}{\sqrt{\kappa(\mathcal{K})}-1}\right)^{2}-1
$$

It follows that $\mathcal{X} \mathcal{F}_{\alpha}^{-1}+\mathcal{F}_{\alpha}^{-1} \mathcal{X}^{T} \succ 0$. Furthermore, it is easy to check that both $\mathcal{X} \mathcal{F}_{\alpha}^{-1} \mathcal{X}^{T}$ and $\mathcal{F}_{\alpha}^{-1}$ are also positive definite. Hence, using 44, we obtain

$$
R(x) \leq \frac{x^{T}\left[\beta \mathcal{X} \mathcal{F}_{\alpha}^{-1} \mathcal{X}+\tau^{2}(1+\alpha) \mathcal{F}_{\alpha}^{-1}\right] x}{x^{T}\left[\beta \mathcal{X} \mathcal{F}_{\alpha}^{-1} \mathcal{X}^{T}+(1+\alpha) \mathcal{F}_{\alpha}^{-1}+\sqrt{\beta(1+\alpha)}\left(\mathcal{X} \mathcal{F}_{\alpha}^{-1}+\mathcal{F}_{\alpha}^{-1} \mathcal{X}^{T}\right)\right] x}<1
$$

from which we deduce that $\lambda_{\text {max }}:=\max R(x)<1$.
The proof of $\lambda_{\min }:=\min R(x) \geq \frac{1}{2(1+\alpha)}$ follows similar arguments as in the second part of the proof of Theorem 3 , with $\mathcal{C}$ and $\mathcal{E}_{\alpha}$ replaced, respectively, by $\mathcal{X}$ and $\mathcal{F}_{\alpha}$.

Remark 1. Note that, using similar arguments as in Theorem 4, we can as well characterize the spectrum of the preconditioned KKT system in the unsteady case, if we define $\mathcal{A}$ as the global coefficient matrix and $\mathcal{P}_{0}$ as

$$
\mathcal{P}_{0}=\left[\begin{array}{cc}
A & 0 \\
0 & \hat{S}
\end{array}\right]
$$

where $A$ and $\hat{S}$ are given by (70) and (73), respectively.
It turns out that, if we specifically use Legendre polynomials and piecewise linear (or bilinear) approximation in the SGFEM discretization of the SOCPs considered herein, then the following result proved by Powell and Elman enables us to further bound the parameter $\alpha$ in Theorem 5 above.

Proposition 2. [25, Lemma 3.7] Let the matrices $G_{k}$ in (27) be defined using normalized Legendre polynomials in uniform random variables on a bounded symmetric interval $[-\nu, \nu]$, and suppose that piecewise linear (or bilinear) approximation is used for the spatial discretization, on quasi-uniform meshes. Let $\left(\lambda_{i}, \varphi_{i}\right)$ be the eigenpairs associated with the $N$-term KLE of the random field $a_{N}$. Then $\kappa(\mathcal{K}) \leq \Phi / \Psi$, where $\Phi=c_{2} \mathbb{E}(a)+\eta$ and $\Psi=c_{1} h^{2} \mathbb{E}(a)-\eta$, with

$$
\eta=c_{2} \sigma_{a} C_{n+1}^{\max } \sum_{i=1}^{N} \sqrt{\lambda_{i}}\left\|\varphi_{i}(\mathbf{x})\right\|_{\infty}
$$

where $C_{n+1}^{\max }$ is the maximal root of the Legendre polynomial of degree $n+1, \sigma_{a}$ is the standard deviation of the random field $a, h$ is the spatial discretization parameter, and $c_{1}$ and $c_{2}$ are constants independent of $h, N$, and $n$.

We can now state the following result.
Corollary 1. Let $\alpha \in[0,+\infty)$. Then, the spectrum of $\hat{S}^{-1} S_{t}$ satisfies

$$
\begin{equation*}
\lambda\left(\hat{S}^{-1} S_{t}\right) \subset\left[\frac{1}{2(1+\alpha)}, 1\right), \quad \forall \alpha<\tilde{\mu}^{2}-1 \tag{76}
\end{equation*}
$$

where $\tilde{\mu}=\frac{1+p+2 \sqrt{p}}{p-1}, p \neq 1$ and $p=\sqrt{\Phi / \Psi}$, with $\Phi$ and $\Psi$ as defined in Proposition 2 .
Proof. The proof is a direct consequence of Proposition 2 and Theorem 5
Next, as the approximation $\hat{S}$ is impractical, we proceed next to derive its practical
version. Now, observe from (64), (65) and (73) that

$$
\begin{align*}
\hat{\mathcal{Z}} & :=\mathcal{K}_{t}+\gamma \mathcal{N} \\
& =\left[\left(I_{n_{t}} \otimes \mathcal{L}\right)+(C \otimes \mathcal{M})\right]+\gamma\left(I_{n_{t}} \otimes \mathcal{M}\right) \\
& =I_{n_{t}} \otimes[\mathcal{L}+\gamma \mathcal{M}]+(C \otimes \mathcal{M}) \\
& =I_{n_{t}} \otimes\left[\left(G_{0} \otimes\left(M+\tau K_{0}\right)+\tau \sum_{i=1}^{N} G_{i} \otimes K_{i}\right)+\gamma\left(G_{0} \otimes M\right)\right]+(C \otimes \mathcal{M}) \\
& =I_{n_{t}} \otimes\left[G_{0} \otimes\left((1+\gamma) M+\tau K_{0}\right)+\tau \sum_{i=1}^{N} G_{i} \otimes K_{i}\right]+(C \otimes \mathcal{M}) \\
& =I_{n_{t}} \otimes\left[G_{0} \otimes \mathcal{Y}+\tau \sum_{i=1}^{N} G_{i} \otimes K_{i}\right]+\left(C \otimes G_{0} \otimes M\right) \tag{77}
\end{align*}
$$

where $\mathcal{Y}=(1+\gamma) M+\tau K_{0}$. Hence, using similar arguments as in Section 2.1 we can now approximate $\hat{\mathcal{Z}}$ using

$$
\begin{equation*}
\hat{\mathcal{Z}}_{0}:=\left(I_{n_{t}} \otimes G_{0} \otimes \mathcal{Y}\right)+\left(C \otimes G_{0} \otimes M\right) \tag{78}
\end{equation*}
$$

In practice, we approximate $\hat{S}$ by applying a multigrid process $n_{t}$ times instead of the inverse of $\mathcal{Y}$ in each of the diagonal blocks of $\hat{\mathcal{Z}}_{0}$ and its transpose.

## 4 Numerical experiments

In this section, we present some numerical results. The numerical experiments were performed on a Linux machine with 80 GB RAM using MATLAB ${ }^{\circledR} 7.14$ together with a MATLAB ${ }^{\circledR}$ version of the AMG code HSL MI20 [4]. We implement each of the mean-based preconditioners $\mathcal{Z}_{0}$ and $\hat{\mathcal{Z}}_{0}$ as given, respectively, by 51 and (78) using one V-cycle of AMG with symmetric Gauss-Seidel (SGS) smoothing to approximately invert $\tilde{K}_{0}$. We remark here that we apply the method as a black-box in each experiment and the set-up of the approximation to $\tilde{K}_{0}$ only needs to be performed once. Unless otherwise stated, in all the simulations, MINRES is terminated when the relative residual error, measured in the Euclidean norm, is reduced to $t o l=10^{-5}$. Note that tol should be chosen such that the truncation tolerance $\epsilon \leq t o l$; otherwise, one would be essentially iterating on the 'noise' from the low-rank truncations, as it were. In particular, we have chosen herein $\epsilon=10^{-8}$.

First, we present our simulations for the cost functional (7) constrained by the steady-state diffusion equation as given by 16 . The random input $a$ is characterized by the covariance function

$$
C_{a}(\mathbf{x}, \mathbf{y})=\sigma_{a}^{2} \exp \left(-\frac{\left|x_{1}-y_{1}\right|}{\ell_{1}}-\frac{\left|x_{2}-y_{2}\right|}{\ell_{2}}\right), \quad \forall(\mathbf{x}, \mathbf{y}) \in[-1,1]^{2}
$$

The forward problem has been extensively studied in, for instance, 25. The eigenpairs $\left(\lambda_{j}, \varphi_{j}\right)$ of the KL expansion of the random field $a$ are given explicitly in [11]. In the
simulations, we set the correlation lengths $\ell_{1}=\ell_{2}=1$ and the mean of the random field $\mathbb{E}[a]=1$.
Next, we investigate the behavior of the solvers (low-rank MINRES and TT-MINRES) for different values of the stochastic discretization parameters $J, P, \sigma_{a}$ as well as $\alpha$ and $\beta$. Moreover, we choose $\xi=\left\{\xi_{1}, \ldots, \xi_{N}\right\}$ such that $\xi_{j} \sim \mathcal{U}[-1,1]$, and $\left\{\psi_{j}\right\}$ are $N$ dimensional Legendre polynomials with support in $[-1,1]^{N}$. The spatial discretization uses $\mathbf{Q}_{1}$ spectral elements. In the considered unsteady SOCP example (that is, in Section (3), the resulting linear systems were solved for time $T=1$.

| LR-Minres | \# iter (t) | \# iter (t) | \# iter (t) | \# iter (t) |
| :--- | :--- | :--- | :--- | :--- |
| $J$ | 481 | 1985 | 8065 | 32513 |
| $P$ |  |  |  |  |
| $\beta=10^{-2}$ |  |  |  |  |
| 20 | $25(32.8)$ | $25(115.4)$ | $27(250.5)$ | $29(736.6)$ |
| 84 | $25(119.7)$ | $27(380.4)$ | $27(582.2)$ | $29(1619.6)$ |
| 210 | $25(141.6)$ | $27(392.8)$ | $27(594.69)$ | $29(1673.9)$ |
| $\beta=10^{-3}$ |  |  |  |  |
| 20 | $21(25.7)$ | $21(113.8)$ | $25(260.9)$ | $25(666.8)$ |
| 84 | $21(128.9)$ | $23(363.7)$ | $25(607.6)$ | $25(1438.1)$ |
| 210 | $21(145.6)$ | $23(385.5)$ | $25(600.8)$ | $25(1471.8)$ |
| $\beta=10^{-4}$ |  |  |  |  |
| 20 | $19(8.2)$ | $21(17.4)$ | $23(67.4)$ | $23(618.3)$ |
| 84 | $19(18.8)$ | $21(42.5)$ | $23(229.7)$ | $23(1313.7)$ |
| 210 | $19(19.6)$ | $21(44.9)$ | $23(576.9)$ | $23(1450.0)$ |
| $\beta=10^{-5}$ |  |  |  |  |
| 20 | $17(19.6)$ | $17(84.8)$ | $21(223.7)$ | $21(578.3)$ |
| 84 | $17(99.9)$ | $19(306.4)$ | $21(520.7)$ | $21(1217.2)$ |
| 210 | $17(115.4)$ | $19(313.63)$ | $21(515.6)$ | $23(1322.6)$ |

Table 1: Results of simulations showing the total number of iterations from lowrank preconditioned MINRES and the total CPU times (in seconds) with $\alpha=1, \beta \in\left\{10^{-2}, 10^{-3}, 10^{-4}, 10^{-5}\right\}, \sigma_{a}=0.1$, and selected spatial $(J)$ and stochastic $(P)$ degrees of freedom.

| $P$ | $J(h)$ | $481\left(\frac{1}{2^{4}}\right)$ | $1985\left(\frac{1}{2^{5}}\right)$ | $8065\left(\frac{1}{2^{6}}\right)$ |
| :--- | :--- | :--- | :--- | :--- |
| $32513\left(\frac{1}{2^{7}}\right)$ |  |  |  |  |
| $20(N=3, n=3)$ | 28,860 | 119,100 | 483,900 | $1,950,780$ |
| $84(N=6, n=3)$ | 121,212 | 500,220 | $2,032,380$ | $8,193,276$ |
| $210(N=6, n=4)$ | 303,030 | $1,250,550$ | $5,080,950$ | $20,483,190$ |

Table 2: Dimension of global coefficient matrix $\mathcal{A}$ in 30 ; here $\operatorname{dim}(\mathcal{A})=3 J P$.
Moreover, our target (or desired state) in both models is the stochastic solution of the
forward model with right hand side 1 and zero Dirichlet boundary conditions ${ }^{1}$
Tables $1,3,4$ and 5 show the results from the low-rank preconditioned MINRES for the model constrained by stationary diffusion. In Table 2 we give the total dimensions of the KKT systems in (30) for various discretization parameters used to obtain the results in Tables 1 . The dimensions range between 28,000 and 20 million.

| LR-Minres | \# iter $(\mathrm{t})$ | \# iter $(\mathrm{t})$ | \# iter $(\mathrm{t})$ |
| :--- | :--- | :--- | :--- |
| $P$ | 20 | 84 | 210 |
| $\operatorname{dim}(\mathcal{A})=3 J P$ | 119,100 | 500,220 | $1,250,550$ |
| $\beta=10^{-3}$ | $19(96.4)$ | $21(336.0)$ | $21(347.93)$ |
| $\beta=10^{-4}$ | $17(86.3)$ | $19(302.6)$ | $19(305.64)$ |
| $\beta=10^{-5}$ | $15(77.4)$ | $17(273.6)$ | $17(283.24)$ |

Table 3: Results of simulations showing the total number of iterations from low-rank preconditioned MINRES and the total CPU times (in seconds) with $\alpha=$ $0, \beta \in\left\{10^{-3}, 10^{-4}, 10^{-5}\right\}, \sigma_{a}=0.1$, and $J=1985\left(h=\frac{1}{2^{5}}\right)$ spatial degrees of freedom, and dimension of the KKT matrix $\mathcal{A}$ solved.

| LR-Minres | \# iter $(\mathrm{t})$ | \# iter $(\mathrm{t})$ | \# iter $(\mathrm{t})$ | \# iter $(\mathrm{t})$ |
| :--- | :--- | :--- | :--- | :--- |
|  | $J$ | $481\left(h=\frac{1}{2^{4}}\right)$ | $1985\left(h=\frac{1}{2^{5}}\right)$ | $8065\left(h=\frac{1}{2^{6}}\right)$ |
|  | $32513\left(h=\frac{1}{2^{7}}\right)$ |  |  |  |
| $\sigma_{a}=0.01$ |  |  |  |  |
| 20 | $17(7.4)$ | $19(16.7)$ | $19(53.4)$ | $21(544.8)$ |
| 84 | $17(17.0)$ | $19(39.0)$ | $19(190.0)$ | $21(1190.0)$ |
| 210 | $17(18.4)$ | $19(40.4)$ | $19(470.0)$ | $21(1230.2)$ |
| $\sigma_{a}=0.4$ |  |  |  |  |
| 20 | $33(13.8)$ | $37(28.0)$ | $41(115.3)$ | $43(1049.8)$ |
| 84 | $35(33.8)$ | $41(84.5)$ | $45(447.0)$ | $47(2610.4)$ |
| 210 | $41(41.9)$ | $47(98.4)$ | $47(782.3)$ | $55(3161.1)$ |

Table 4: Results showing the effect of significant decrease ( $\sigma_{a}=0.01$ ) and increase ( $\sigma_{a}=0.4$ ) in the standard deviation of the random input $a$ on the low-rank preconditioned MINRES with $\alpha=1, \beta=10^{-4}$ and selected spatial $(J)$ and stochastic $(P)$ degrees of freedom.

We have solved the linear systems using our proposed block-diagonal preconditioner, together with the approximation $S_{1}$ for the Schur complement as given by (41) and (51). We observe first, from Table 1, that our preconditioner is robust with respect to the discretization parameters. Herein, $h$ is the spatial mesh size. Moreover, as

[^0]illustrated in Table 4 it performs relatively better if the standard deviation $\sigma_{a}$ of the random input is smaller than when it is increased. Nonetheless, the preconditioner still maintains its robustness in $\{0.01,0.4\}$. For $\sigma_{a}>0.5$, we can no longer guarantee the positive-definiteness of the matrix $\mathcal{K}$ corresponding to the forward problem [25].

The results in Tables 1. 4, and 5 were obtained with $\alpha=1$, whereas those in a Table 3 were computed with $\alpha=0$. We have reported in Table 5 the values of the tracking term and the cost functional for $\alpha=1$ and $\sigma_{a}=0.1$. As expected, the tracking term gets smaller and smaller as the regularization parameter $\beta$ decreases, and the cost functional also decreases accordingly converging, respectively, to $1.2 \times 10^{-4}$ and $2.5 \times 10^{-4}$.

| $\beta$ | $10^{-2}$ | $10^{-4}$ | $10^{-6}$ | $10^{-10}$ |
| :--- | :--- | :--- | :--- | :--- |
| $\\|y-\bar{y}\\|_{L^{2}(\mathcal{D}) \otimes L_{\rho}^{2}(\Gamma)}^{2}$ | $5.1 \times 10^{-3}$ | $1.8 \times 10^{-4}$ | $1.2 \times 10^{-4}$ | $1.2 \times 10^{-4}$ |
| $\mathcal{J}(y, u)$ | $1.4 \times 10^{-2}$ | $4.2 \times 10^{-4}$ | $2.5 \times 10^{-4}$ | $2.5 \times 10^{-4}$ |

Table 5: Tracking term and the cost functional in the steady-state model for different values of $\beta$ and with $\alpha=1, \sigma_{a}=0.1, J=1985\left(h=\frac{1}{2^{5}}\right), P=84(N=$ $6, n=3)$.


Figure 1: The mean (left) and variance (right) of the target using low-rank preconditioned MINRES for the steady-state model with $\sigma_{a}=0.1, J=1985$ ( $h=$ $\left.\frac{1}{2^{5}}\right), P=84(N=6, n=3)$.

Figure 1 shows the mean and the variance of the target function, whereas Figures 2 and 3, respectively, depict the statistics for the state and control functions computed with $\alpha=1$. Here, we see that the mean of the state visually coincides with the mean of the desired state. However, as anticipated, the variance of the state is reduced. When
we set $\alpha=0$, we observe that the solver recovers both the mean and variance of the state as illustrated in Figure 4 This observation, together with the results in Table 3 , further confirms the robustness of our preconditioner for $\alpha \in\{0,1\}$.


Figure 2: The mean (left) and variance (right) of the state using low-rank preconditioned MINRES for the steady-state model with $\alpha=1, \beta=10^{-2}, \sigma_{a}=$ $0.1, J=1985\left(h=\frac{1}{2^{5}}\right), P=84(N=6, n=3)$ and truncation tolerance $\epsilon=$ $10^{-8}$.

Next, we present in Table 6 our simulation results for the unsteady diffusion constrained model as discussed in Section [3. Here, for $\alpha \in\{0,1\}$ and different values of $\beta$, we show the outputs of our simulations showing the total CPU times and the total number of iterations from preconditioned MINRES in TT-format. Also, $\mathrm{DoF}=J \cdot P \cdot n_{t}$ is the size of each of the 9 block matrices in $\mathcal{A}$. That is, the optimality matrix $\mathcal{A}$ is of dimension 3 DoF . In particular, we have done the computations with $J=1985\left(h=\frac{1}{2^{5}}\right), P=56(N=5, n=3), \sigma_{a}=0.1$, and different numbers of total time steps $n_{t}$.

As in the steady-state case, we see from Table 6 that TT-MINRES, when used together with our mean-based preconditioner as given by $(72)$ and $(78)$ is quite robust, but in general yields fewer iterations for $\alpha=0$ than for $\alpha=1$. We remark here that we used a smaller tolerance tol $=10^{-3}$ in the unsteady case because MATLAB ${ }^{\circledR}$ took more time because of the rapid the growth of TT-ranks. Although not reported here, we also got robust two-digit TT-MINRES iterations when we used $t o l=10^{-5}$.


Figure 3: The mean (left) and variance (right) of the control using low-rank preconditioned MINRES for the steady-state model with $\alpha=1, \beta=10^{-2}, \sigma_{a}=$ $0.1, J=1985\left(h=\frac{1}{2^{5}}\right), P=84(N=6, n=3)$ and truncation tolerance $\epsilon=$ $10^{-8}$.

| TT-Minres | \# iter $(\mathrm{t})$ | \# iter $(\mathrm{t})$ | \# iter $(\mathrm{t})$ |  |  |  |
| :--- | :--- | :--- | :--- | :---: | :---: | :---: |
| $n_{t}$ | $2^{5}$ | $2^{6}$ | $2^{8}$ |  |  |  |
| $\operatorname{dim}(\mathcal{A})=3 J P n_{t}$ | $10,671,360$ | $21,342,720$ | $85,370,880$ |  |  |  |
| $\alpha=1$, tol $=10^{-3}$ |  |  |  |  |  |  |
| $\beta=10^{-5}$ | $6(285.5)$ | $6(300.0)$ | $8(372.2)$ |  |  |  |
| $\beta=10^{-6}$ | $4(77.6)$ | $4(130.9)$ | $4(126.7)$ |  |  |  |
| $\beta=10^{-8}$ | $4(56.7)$ | $4(59.4)$ | $4(64.9)$ |  |  |  |
| $\alpha=0$, tol $=10^{-3}$ |  |  |  |  |  |  |
| $\beta=10^{-5}$ | $4(207.3)$ | $6(366.5)$ | $6(229.5)$ |  |  |  |
| $\beta=10^{-6}$ | $4(153.9)$ | $4(158.3)$ | $4(172.0)$ |  |  |  |
| $\beta=10^{-8}$ | $2(35.2)$ | $2(37.8)$ | $2(40.0)$ |  |  |  |

Table 6: Results of simulations of the model with time-dependent diffusion constraint showing the total number of iterations from preconditioned MINRES in TTformat and the total CPU times (in seconds) for selected parameter values and degrees of freedom.

## 5 Conclusions and outlook

In this paper, we have derived and implemented block-diagonal Schur complementbased preconditioners for linear systems arising from the SGFEM discretization of the


Figure 4: The mean (left) and variance (right) of the state using low-rank preconditioned MINRES for the steady-state model with $\alpha=0, \beta=10^{-4}, \sigma_{a}=$ $0.1, J=1985\left(h=\frac{1}{2^{5}}\right), P=84(N=6, n=3)$ and truncation tolerance $\epsilon=$ $10^{-8}$.
stochastic optimal control problems constrained by either stationary or time-dependent PDEs with random inputs. Moreover, we analyzed the spectra of the derived preconditioners. Our approach to the solution of the KKT linear systems entails a formulation that solves the systems at once (for all time steps in the unsteady case). This strategy leads, more often than not, to a large system that cannot be solved using direct solvers. However, combining our proposed preconditioners with the low-rank solvers considered herein have proven efficient in accomplishing the tasks.
Although the TT-MINRES works quite well for the time-dependent problem considered in this paper, the rapid growth of the TT-ranks is not a trivial issue. In a related work, we are currently exploiting some capabilities of the TT toolbox to minimize the rank growth and hence make the solver a lot more efficient.

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[^0]:    ${ }^{1}$ Note that this is not an 'inverse crime' as the right-hand side of the forward model used is deterministic, unlike in the state equation.

